Discrete-Event Simulation: Multi-Stream Lehmer RNGs

Lawrence M. Leemis and Stephen K. Park, Discrete-Event Simulation - A First Course, Prentice Hall, 2006

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Discussion on Project and Other Assignments

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- ▶ Project 1

Introduction

- Typical DES models have many stochastic components.
 - e.g., arrivals and services
- Have a unique randomness for each stochastic component
- One option: multiple RNGs
 - Often considered a poor option
- ▶ One RNG with multiple "streams" of random numbers
 - One stream per stochastic component
 - Considered a better option
 - Method: partition output from a RNG into multiple streams
 - We have been using the Lehmer RNG from the authors of the textbook

Lehmer RNG

Lehmer RNG used in ssq2 and sis2

```
▶ in C/C++: rng.h and rng.c
  double Random(void);
  void PutSeed(long x);
 void GetSeed(long *x);
  void TestRandom(void);
▶ in Java (and C++): Rng.java (and Rng.cpp with minor syntax
 difference)
  class Rng {
      [\ldots]
      public double random() {[...]}
      public void putSeed(long x) {[...]}
      public long getSeed() {[...]}
      public void testRandom() {[...]}
```

Revisit Simulation Program ssq2

 ssq2 has two stochastic components: arrival process and service process, e.g.,

```
double GetArrival(void) {
    static double arrival = START;

    arrival += Exponential(2.0);
    return (arrival);
}

double GetService(void) {
    return (Uniform(1.0, 2.0));
}
```

Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- First attempt: partition output from our Lehmer RNG into multiple streams
 - Method: allocate and retain internal state of the RNG for each stochastic process

Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- Allocate a different generator state variable to each process and retain it before switching to the other process
 - ▶ In the Lehmer RNG, the generator state is fully represented by the seed
 - Allocate to the service process its own static variable and initialized with a value

GetService with Unique Seed

```
double GetService(void) {
    double s;
    static long x = 12345;
    PutSeed(x);
    s = Uniform(1.0, 2.0);
    GetSeed(&x);
    return (s);
}
```

Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

► Allocate to arrival process its own *static* variable and initialized with a *different* value from the service process

```
GetArrival with Unique Seed
double GetArrival(void) {
    static double arrival = START;
    static long x = 54321;
    PutSeed(x);
    arrival += Exponential(2.0);
    GetSeed(&x);
    return (arrival);
```

where x represents the current state of the service process

Modified Arrival and Service Processes

- As modified, arrival and service times are drawn from different streams of random numbers
- Provided the streams do not overlap, the processes are uncoupled
- ▶ Although the choice of seed for each stream is deceivingly simple, the choices may in fact be *poor* ones.
- ► Execution time cost is negligible (see Example 3.2.3 in next slide)

Exericse L6-1

Complete Example 3.2.3 as instructed as follows,

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2b*
- Replace GetArrival and GetService by those introduced in the 1st approach
- ▶ Set *LAST* to 1,000,000 jobs
- Compile and run the program. Measure the execution time of new and original programs
- ► Compare the results with from those using *ssq2* (without any modification). Are they the same?
- ► Compare execution time of the new and original programs, how much slower is the new program?

Exericse L6-1: Hints

To measure program execution time in Windows, use *PowerShell*. Below is an example that measures the execution time of a *dir* command,

Measure Program Execution Time

```
C:> powershell
PS C:> Measure-Command { dir | Out-Default }
  Directory: C:\
 Mode
                     LastWriteTime
                                      Length Name
  [...]
  Days
  Hours
  Minutes
  Seconds
  Milliseconds
                   : 129
                   : 1299779
  Ticks
  TotalDays
                   : 1.50437384259259E-06
  TotalHours : 3.6104972222222E-05
  TotalMinutes : 0.00216629833333333
  TotalSeconds : 0.1299779
  Total Milliseconds: 129.9779
```

See

https://technet.microsoft.com/en-us/library/hh849910.aspx

Exercise L6-1: Hints

To measure program execution time in Linux, use the *time* command. Below is an example that measures the execution time of a *ls* command,

Measure Program Execution Time

For more detail on the command, see its manual page (i.e., man time).

Streams using Multiple Seeds: Discussion

- Objective: allocate a unique stream of random numbers of each stochastic component
 - Examples of stochastic components: arrival and service processes
- ▶ Discussed approach: using multiple seeds of RNGs to produce multiple unique streams of random numbers
- ▶ Potential problem: assignment of initial seeds (or initial state)
 - ▶ Initial states should be chosen to produce *disjoint* streams
 - ▶ If states are picked at whim, no guarantee of disjoint streams
 - Some initial states could be just a few calls to Random() away from one another

Jump Multipliers

- ▶ Objective: produce multiple *disjoint* streams of random numbers
- ▶ Theorem 3.2.1 is the key to creating streams

Theorem 3.2.1

Given $g(x) = ax \mod m$ and integer j with j = 1, 2, ..., m - 1, the associated *jump function* is

$$g^{j}(x) = (a^{j} \mod m)x \mod m \tag{1}$$

and has the jump multiplier $a^j \mod m$ if $(g(\cdot))$ generates x_0, x_1, x_2, \ldots then $g^j(\cdot)$ generates x_0, x_j, x_{2j}, \ldots

Numerical Examples

- If m = 31 and a = 3 and j = 6, the jump multiplier is $a^j \mod m = 36 \mod 31 = 16$
- ▶ If $x_0 = 1$ then $g(x) = 3x \mod 31$ generates

$$\underline{1}$$
, 3, 9, 27, 19, 26, $\underline{16}$, 17, 20, 29, 25, 13, $\underline{8}$, 24, 10, 30, 28, 22, $\underline{4}$, 12, 5, 15, 14, 11, $\underline{2}$, 6 . . .

► The jump function $g^6(x) = 16x \mod 31$ generates $1, 16, 8, 4, 2, \ldots$ i.e., the first sequence is x_0, x_1, x_2, \ldots ; the second is x_0, x_6, x_{12}, \ldots

Program for Numerical Example

```
#include < stdio.h>
#include <math.h>
int main() {
    long m = 31, a = 3, j = 6, jm, i, n =
          32. \times = 1:
    jm = (long)pow(a, j) \% m;
    printf("%2Id", x);
    for (i = 0: i < n: i ++) {
         printf(", %2Id", x = a * x % m);
    printf("\n\nJump\ Multiplier = \%ld\n\n".
         im);
    x = 1:
    printf("%2Id", x);
    for (i = 0; i < n/6; i ++) {
         printf(", \%2Id", x = im * x \% m):
    printf("\n");
    return 0:
```

```
import iava.io.*:
import iava.lang.Math:
import java.text.*;
public class Ex3_2_4{
    public static void main(String[] args) {
        String format = "%2d":
        long m = 31, a = 3, j = 6, jm, i, n
             = 32. x = 1:
       im = (long) Math.pow(a, j) % m;
        System.out.format(format, x);
        for (i = 0; i < n; i ++) {
            System.out.format(", " +
                 format x = a * x % m:
        System.out.format("\n\nJump
             Multiplier = " + format +
             "\n\n", jm);
        x = 1; System.out.format(format, x);
        for (i = 0; i < n/6; i ++) {
            System.out.format(", " +
                 format, x = im * x % m);
        System.out.print("\n");
```

Using Jumper Function

- 1. Compute the jump multiplier $g^j(\cdot) = a^j \mod m$, which is a one time cost.
- 2. $g^{j}(\cdot)$ permits jumping from x_0 to x_j to x_{2j} to ...
- 3. User supplies one initial seed
- 4. If j is chosen well, $g^{j}(\cdot)$ can "plant" additional initial seeds
- 5. Each planted seed corresponds to a different stream
- 6. Each planted seed is separated by *j* calls to *Random()*

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Maximal Modulus-Compatible Jump Multiplier

Definition 3.2.1

Given a Lehmer random-number generator with prime modulus m, full-period modulus-compatible multiplier a, and a requirement for s disjoint streams as widely separated as possible, the maximal jump multiplier is $a^j \mod m$, where j is the largest integer less than $\lfloor m/s \rfloor$ such that $a^j \mod m$ is modulus-compatible with m.

Example 3.2.6

Jump multipliers for $(a, m) = (48271, 2^{31} - 1) = (48271, 2147483647)$ RNG			
# of streams s	$\lfloor m/s \rfloor$	jump size <i>j</i>	
1024	$\lfloor \frac{2^{31}-1}{1024} \rfloor =$	2082675	48271 ²⁰⁸²⁶⁷⁵
	2097151		$mod\ 2147483647 = 97070$
512	$\lfloor \frac{2^{31}-1}{512} \rfloor =$	4170283	48271 ⁴¹⁷⁰²⁸³
	4194303		$mod\ 2147483647 = 44857$
256	$\lfloor \frac{2^{31}-1}{256} \rfloor =$	8367782	48271 ⁸³⁶⁷⁷⁸²
	8388607		$mod\ 2147483647 = 22925$
128	$\left\lfloor \frac{2^{31}-1}{128} \right\rfloor =$	16775552	48271 ¹⁶⁷⁷⁵⁵⁵²
	16777215		$mod\ 2147483647 = 40509$

Jump Multiplier Search Program in C

Definition 2.2.1

The multiplier a is modulus-compatible with the prime modulus m if and only if r < qwhere $r = m \mod a$ and $q = \lfloor m/a \rfloor$

Functions in Jump Multiplier Search Program

```
long modular-pow(long long base, long long exponent, long long modulus);
long find_m_compatible(long upper, long a, long m)
    long i, n, r, q;
    for (i = upper; i >= 1; i --) {
        n = modular_pow((long long)a, (long long)i, (long long)m);
        r = m \% n:
       q = m / n;
        if (r < q) { /* if n is modulus-compatibl with m */
            return i:
    return 0;
```

This function is to compute base exponent mod modulus

Functions in Jump Multiplier Search Program

```
* Reference
* http://www.sanfoundry.com/cpp-program-implement-modular-exponentiation-algorithm/
* */
long modular_pow(long long base, long long exponent, long long modulus)
    long long result = 111;
    while (exponent > 0 II) {
        if (exponent % 2|| == 1||) {
            result = (result * base) % modulus;
        exponent = exponent >> 1 | 1 |;
        base = (base * base) % modulus:
    return (long) result;
```

Jump Multiplier Search Program in Java

Jump Multiplier Search Program in Java

```
. . . . . . . . .
 private long modular_pow(long base, long exponent, long modulus) {
     long result = 11;
     while (exponent > 01) {
         if (exponent % 2| == 1|) {
             result = (result * base) % modulus;
         exponent = exponent >> 11;
         base = (base * base) % modulus;
     return result;
```

Jump Multiplier Search Program in Java

```
private long find_m_compatible(long upper, long a, long m) {
    long i, n, r, q;
    for (i = upper; i >= 1; i --) {
        n = modular_pow(a, i, m):
        r = m \% n:
        q = m / n;
        if (r < q) { /* if n is modulus-compatibl with m */
            return i:
    return 0;
```

Exericse L6-2

- ▶ Use the program discussed above to compute the jump multiplier table similar to slide 19.
 - ▶ You will enter, compile, and run the programs.

Library rngs

- lacktriangledown rngs is an upward-compatible multi-stream replacement for \emph{rng}
- ▶ By default, provides 256 streams, indexed 0 to 255 (0 is the default)
- Only one stream is active at any time
- Six available functions:
 - Random(void)
 - PutSeed(long x): superseded by PlantSeeds
 - GetSeed(long *x)
 - ► TestRandom(void)
 - SelectStream(int s): used to define the active stream
 - ▶ PlantSeeds(long x): "plants" one seed per stream
- ▶ Henceforth, *rngs* is the library of choice

Example 3.2.7: ssq2 Revisited

- Use rngs functions for GetArrival, GetService
- ▶ Include *rngs.h* and use *PlantSeeds*(12345)

GetArrival Method

```
double GetArrival(void) {
    static double arrival = START;
    SelectStream(0);
    arrival += Exponential(2.0);
    return (arrival);
}
```

GetService Method

```
double GetService(void) {
    SelectStream(2);
    return (Uniform(1.0, 2.0));
}
```

Exericse L6-3

Complete Example 3.2.7 as instructed as follows,

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2c*
- Replace GetArrival and GetService by those using library rngs. Find library rngs in Blackboard.
- ▶ Set *LAST* to 1,000,000 jobs
- Compile and run the program. Measure the execution time of new and original programs
- ► Compare the results with from those using *ssq2* (without any modification) and those using *ssq2b*(Exercise L6-1). Are they the same?
- ► Compare execution time of the new and originial programs, how much slower is the new program?

Uncoupling Stochastic Processes

- Per modifications, arrival and service processes are uncoupled
- Consider changing the service process to Uniform(0.0, 1.5) + Uniform(0.0, 1.5)
- Without uncoupling, arrival process sequence would change!
- With uncoupling, the service process sees exactly the same arrival sequence
- ► Important variance reduction technique

Single-Server Service Node with Multiple Job Types

- Extend the single-server service node model from Chapter 1
- Consider multiple job types, each with its own arrival and service process
- ► Examples 3.2.8 and 3.2.9: Suppose there are two job types
 - 1. Exponential (4.0) interarrivals, Uniform (1.0, 3.0) service
 - 2. Exponential(6.0) interarrivals, Uniform(0.0, 4.0) service

Use rngs to allocate a different stream to each stochastic process

Arrival Process for Multiple Job Types

The arrival process generator in program ssq2 can be modified as follows,

Example 3.2.8: Arrival Process

```
double GetArrival(int *j) { /* returns job type in j */ const double mean[2] = {4.0, 6.0}; /* two job types */
    static double arrival[2] = {START, START};
    static int init = 1:
    double temp;
                                            /* initialize the arrival arrav */
    if (init) {
        SelectStream (0):
         arrival[0] += Exponential(mean[0]);
        SelectStream (1);
         arrival[1] += Exponential(mean[1]):
         init = 0:
    if (arrival[0] <= arrival[1])
        *i = 0:
                                            /* next arrival is job type 0 */
    else
                                            /* next arrival is job type 1 */
        *i = 1:
    temp = arrival[*j];
                                           /* next arrival time to be returned */
    SelectStream(*i);
    arrival[*i] += Exponential(mean[*i]); /* arrival after next arrival */
    return (temp);
```

Service Process for Multiple Job Types

The service process generator in program ssq2 can be modified as follows,

Example 3.2.9: Service Progress

```
double GetService(int j)
{
    const double min[2] = {1.0, 0.0};
    const double max[2] = {3.0, 4.0};

/*
    * Two RNG streams, i.e., streams 0 and 1 are used in the arrival
    * process generator. We now use streams 2 and 3 for the service
    * process generator. In the following, j should be either 0 or 1.
    * */
    SelectStream(j + 2);
    return (Uniform(min[j], max[j]));
}
```

Service Process for Multiple Job Types

- ▶ Index *j* matches service time to appropriate job type
- ▶ All four simulated stochastic processes are uncoupled
- ► Any process could be changed without altering the random sequence of others!

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Consistency Check

- Additional modification to ssq2
 - job-type-specific statistics-gathering needs to be added in the main method/function
- ▶ How do we know if our modifications are correct? *Use consistency check to increase confidence.*
 - $\overline{w} = \overline{d} + \overline{s}$
 - $ightharpoonup \overline{I} = \overline{q} + \overline{x}$
 - ► How about average service time of both job types? Since (1.0 + 3.0)/2 = (0.0 + 4.0)/2 = 2.0, we expect $\overline{s} = 2.0$
 - ► How about the net arrival rate of both job types? Since the arrival rates of job types 0 and 1 are 1/4 and 1/6, respectively, we expect the net arrival rate should be 1/4 + 1/6 = 5/12, i.e., $\overline{r} = 1/(5/12) = 12/5 = 2.4$.
 - ► The steady-state utilization should be the ratio of the arrival rate to the service rate, i.e., $(5/12)/(1/20 = 5/6 \approx 0.83$

Exercise L6-4

Modify program *ssq2* to support two job types as discussed.

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2d*
- ▶ Modify program *ssq2d* as suggested in Examples 3.2.8 and 3.2.9.
- Modify the main method/function to include job-type-specific statisics
- Answer the following questions,
 - What portion of processed jobs are of type 0?
 - ▶ What are \overline{w} , \overline{d} , \overline{s} \overline{l} , \overline{q} , and \overline{x} for each job type?
 - What did you do to convince yourself that your results are valid (hint: consistency check)?
 - ▶ Why are \overline{w} , \overline{d} , and \overline{s} the same for both job types, while \overline{l} , \overline{q} , and \overline{s} are different?

Summary

- Multiple stream RNGs
 - Generate disjoint random number streams
 - Want the streams are far apart
- More simulation examples to be discussed