Discrete-Event Simulation

Lawrence M. Leemis and Stephen K. Park, Discrete-Event Simul A First Course, Prentice Hall, 2006

Hui Chen

Computer Science Virginia State University Petersburg, Virginia

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Introduction

- Programs ssq1 and sis1 are trace-driven discrete-event simulations
 - Both rely on input data from an external source
- These realizations of naturally occurring stochastic processes are limited
- Cannot perform "what if" studies without modifying the data
- Solution
 - Convert the single server service node and the simple inventory system to utilize *randomly generated input*
 - Use a random-number generator to produce the randomly generated input
 - Discrete-event simulation programs using the randomly generated input does not depend on external trace data

Single Queue Service Node: Revisited

- Need two stochastic assumptions
 - arrival times
 - service times
- The assumptions governs how arrival and service times are randomly generated in discrete-event simulation programs

- Service time
 - Range: between 1.0 and 2.0
 - Distribution within the range? Without further knowledge, we assume no time is more likely than any other
 - To generate service times: use u = Uniform(1.0, 2.0) random variate

Is it reasonable to assume that service times are uniformly distributed, e.g., service times are generated using u = Uniform(1.0, 2.0) random variate?

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It depends.

In most applications, it is unrealistic to assume service times are uniformly distributed.

Service Time in *ssq1.dat* Trace Data

Is service times in *ssq1.dat* uniformly distributed?



Example: Generating Service Times: Exponential Distribution

- In general, it is unreasonable to assume that all possible values are equally likely.
- Frequently, small values are more likely than large values
- ▶ Need a non-linear transformation that maps $0 \rightarrow 1$ to $0 \rightarrow \infty$ since 0 < u = Uniform(0, 1) < 1

Single Queue Service Node Exponential Random Variate

Example: Generating Service Times: Exponential Distribution

A common nonlinear transformation is

$$x = -\mu \ln(1 - u) \tag{1}$$

The transformation is monotone increasing, one-to-one, and onto

$$0 < \mu < 1 \iff 0 > -u > -1$$
 (2)

$$\iff 0+1 > -u+1 < -1+1 \tag{3}$$

$$\iff 1 > 1 - u > 0 \tag{4}$$

$$\iff \ln(1) > \ln(1-u) > \ln(0) \tag{5}$$

$$\iff 0 > ln(1-u) > \infty$$
 (6)

$$\iff 0 < -\ln(1-u) < \infty \tag{7}$$

$$\iff 0 < -\mu \ln(1-u) < \infty \tag{8}$$

$$\iff 0 < x < \infty \tag{9}$$

Example: Generating Service Times: Exponential Distribution

► The common nonlinear transformation x = -µln(1 - u) is monotone increasing, one-to-one, and onto

$$0 < \mu < 1 \iff 0 < -\mu \ln(1-u) < \infty \iff 0 < x < \infty$$
 (10)

which generates $Exponential(\mu)$ random variate



Figure: Exponential-variate-generation Geometry

Single Queue Service Node Exponential Random Variate

Example: Generating Service Times: Exponential Distribution

The common nonlinear transformation

$$x = -\mu \ln(1 - u) \tag{11}$$

generates *Exponential*(μ) random variate • Note that 0 < u < 1 and

$$\int_{0}^{1} -\mu \ln(1-u) du = -\mu \int_{0}^{1} \ln(1-u) du$$
(12)

$$= -\mu \int_0^1 -\ln(1-u)d(1-u) = \mu \int_0^1 \ln(1-u)d(1-u)$$
(13)

$$= \mu \{ \ln(1-u)(1-u)|_{0}^{1} - \int_{0}^{1} (1-u) d\ln(1-u) \}$$
(14)

$$= \mu \{ 0 - (1 - u) \frac{1}{1 - u} (1 - u) |_0^1 \}$$
(15)

$$= -\mu(1-u)\frac{1}{1-u}(1-u)|_{0}^{1} = -\mu(1-u)|_{0}^{1}$$
(16)

$$= \mu$$
 (17)

i.e., the parameter μ specifies the sample mean

Generating *Exponential*(μ) Random Variate

Definition 3.1.1 ANSI C Function for *Exponential*(μ)

```
double Exponential(double \mu)
```

```
return - \mu * log(1.0 - Random());
```

where Random() generates u = Uniform(0, 1) random variate and μ is the sample mean.

Example: Generating Service Times: Exponential Distribution

In the single-server service node simulation, we use Exponential(μ_s) to generate service times,

$$s_i = Exponential(\mu_s); \quad i = 1, 2, 3, \dots, n$$
 (18)

where μ_s is the sample mean of service times.

Example: Generating Interarrival Times: Exponential Distribution

In the single-server service node simulation, we use $Exponential(\mu_a)$ to generate interarrival times,

$$a_i = a_{i-1} + Exponential(\mu_a); \quad i = 1, 2, 3, \dots, n$$
 (19)

where μ_a is the sample mean of interarrival times.

Example: Recap

- Inter-arrival times
 - Generating u = Uniform(a, b) random variate
 - ► Generating *u* = *Exponential*(*a*) random variate
- Service times
 - Generating u = Uniform(a, b) random variate
 - Generating u = Exponential(a) random variate

Simulation Program ssq2

- Program ssq2 is an extension of ssq1
 - Interarrival times are drawn from Exponential(2.0)
 - Service times are drawn from Uniform(1.0, 2.0)
- The program generates job-averaged and time-averaged statistics
 - ► *T*: average interarrival time
 - w: average wait
 - ▶ d: average delay
 - ► <u>s</u>: average service time
 - \overline{I} : average # in the node
 - \overline{q} : average # in the queue
 - x: server utilization

Exercise L4-1

In this exercise, you are required to complete the following tasks,

- Develop ssq2 by revising ssq1 program.
- Compile and run the ssq2 program.
- When writing the program, meet the following,
 - Interarrival times are drawn from Uniform(0.0, 6.0)
 - Service times are drawn from Exponential(2.0)
- Submission: the source code of *ssq2*, the results of the program, and evidence that your program appears to be correct.

Example 3.1.3: Theoretical Result from Analytic Model

 The theoretical averages for a single-server service node using *Exponential*(2.0) inter-arrivals and *Uniform*(1.0, 2.0) service times are (Gross and Harris, 1985),

$$\overline{r}$$
 \overline{w} \overline{d} \overline{s} \overline{l} \overline{q} \overline{x}
2.00 3.83 2.33 1.50 1.92 1.17 0.75

- Although the server is busy only 75% of the time, on average there are approximately two jobs in the service node
- A job can expect to spend more time in the queue than in service
- ► To achieve these averages, many jobs must pass through node

Example 3.1.3: Results from Simulation Program *ssq2*

The accumulated average wait was printed every 20 jobs



Figure: Average wait times

► The convergence of *w* is slow, erratic, and dependent on the initial seed

Use of Program ssq2

- The program can be used to study the steady-state behavior
 - Will the statistics converge independent of the initial seed?
 - How many jobs does it take to achieve steady-state behavior?
- It can be used to study the transient behavior
 - Fix the number of jobs processed and replicate the program with the initial state fixed
 - Each replication uses a different initial rng seed

Exericse L4-2

You are required to reproduce the figure in slide 20. You may take steps below (using the C/C++ program as an example),

- Convert the main function int main(void) to function void SimulateOnce(long seed, long last).
 - seed: seed of RNG
 - *last*: the number of jobs to process
- Add the main function in which you call SimulateOnce with seed and last in a loop with last as the loop variable to simulate with the number of jobs as 20, 40, ..., 1000.
- ► Format the output in the "CSV" format.
- Run the program and graph the results.
- Submission: program source code, running results, and graph.

Geometric Random Variables

The Geometric(p) random variate is the discrete analog to a continuous Exponential(µ) random variate Let x = Exponential(µ) = µln(1 − µ), y = |x|, and p = Pr(y ≠ 0)

$$y = \lfloor x \rfloor \neq 0 \Longleftrightarrow x \ge 1 \tag{20}$$

$$\iff \mu \ln(1-\mu) \ge 1$$
 (21)

$$\iff \ln(1-\mu) \le -1/\mu \tag{22}$$

$$\iff 1 - \mu \le e^{-1/\mu} \tag{23}$$

Since $1 - \mu$ is also Uniform(0.0, 1.0) and $p = Pr(y \neq 0) = e^{-1/\mu}$ Finally, since $\mu = -1/ln(p)$, $y = \lfloor ln(1-\mu)/ln(p) \rfloor$

Generating Geometric(p) Random Variates

Definition 3.1.2 ANSI C Function for Geometric(p)

```
long Geometric(double p) use 0.0
```

```
return (long)(log(1.0 - Random()) / log(p));
```

- Random() generates u = Uniform(0,1) random variate.
- The mean of a Geometric(p) random variate is p/(1-p)
 - If p is close to zero then the mean will be close to zero
 - If p is close to one, then the mean will be large

ł

Example 3.1.4: Composite Service Model

Now consider a composite service model

- Assume that jobs arrive at random with a steady-state arrival rate of 0.5 jobs per minute
- Assume that Job service times are composite with two components
 - ► The number of service tasks is 1 + *Geometric*(0.9)
 - ▶ The time (in minutes) per task is *Uniform*(0.1, 0.2)

Example 3.1.4: Composite Service Model

ANSI C Function for the Composite Service Model

```
double GetService(void)
```

```
long k;
double sum = 0.0;
long tasks = 1 + Geometric(0.9);
for (k = 0; k < tasks; k++)
sum += Uniform(0.1, 0.2);
return (sum);
```

Example 3.1.4: Composite Service Model: Analytic Model

The theoretical steady-state statistics for this model are

$$\overline{r}$$
 \overline{w} \overline{d} \overline{s} \overline{l} \overline{q} \overline{x}
2.00 5.77 4.27 1.50 2.89 2.14 0.75

- The arrival rate, service rate, and utilization are identical to Example 3.1.3 (See slide 19)
- The other four statistics are significantly larger
- Performance measures are sensitive to the choice of service time distribution

Simple Inventory System: Example 3.1.5

- Program sis2 has randomly generated demands using an Equilikely(a, b) random variate
- Using random data, we can study transient and steady-state behaviors
- If (a, b) = (10, 50) and (s, S) = (20, 80), then the approximate steady-state statistics are

$$\overline{d}$$
 $\overline{\sigma}$ \overline{u} \overline{l}^+ \overline{l}^-
30.00 30.00 0.39 42.86 0.26

Exercise L4-3

In this exercise, you are required to complete the following tasks,

- Compile and run the *sis2* program. Document the results.
- ▶ Make a copy of the *sis2* program, revise it to meet the following,
 - The demand is drawn from Geometric(0.967742)

and then compile and run the program.

Submit your work including both version of the sis2 program and the results of both runs in Blackboard

Effects of Number of Time Intervals and Seed of RNG

The average inventory level \$\overline{l} = \overline{l}^+ - \overline{l}\$ approaches steady state after several hundred time intervals



Figure: Number of Time Intervals (n)

Convergence is slow, erratic, and dependent on the initial seed

Exercise L4-4

You are required to reproduce the figure in slide 30. You may take steps below (using the Java program as an example),

- Convert the main function public static void main(String[] args) to function public static void SimulateOnce(long seed, long stop).
 - seed: seed of RNG; stop: the number of intervals to process
 - Format the output in the "CSV" format
- Add the public static void main(String[] args function in which you call SimulateOnce with seed and stop in a loop with stop as the loop variable to simulate with the number of intervals as 5, 10, 15, ..., 200.
- Run the program and graph the results
- Submission: program source code, results, and graph.

Simple Inventory System

Example 3.1.7: Optimal Inventory Policy

▶ If we fix *S*, we can find the optimal cost by varying *s*



Figure: Dependent Cost for (s, S) Inventory System

where $c_{setup} = \$1,000$, $c_{hold} = 25$, $c_{short} = 700$, min(DependentCost) = \$1,624.86, and s = 24.

Recall that the dependent cost ignores the fixed cost of each item

Example 3.1.7: Discussion

- Using a fixed initial seed guarantees the exact same demand sequence
 - Any changes to the system are caused solely by the change of s
- A steady state study of this system is unreasonable
 - All parameters would have to remain fixed for many years
 - When n = 100 we simulate approximately 2 years
 - When n = 10000 we simulate approximately 192 years

Statistical Considerations

- Example 3.1.7 illustrates two consideration
 - Variance reduction
 - Robust estimation
- With Variance Reduction, we eliminate all sources of variance except one
 - > Transient behavior will always have some inherent uncertainty
 - We kept the same initial seed and changed only s
- Robust Estimation occurs when a data point that is not sensitive to small changes in assumptions
 - Values of s close to 23 have essentially the same cost
 - Would the cost be more sensitive to changes in S or other assumed values?

Exercise L4-5

You are required to reproduce the figure in slide 32. Hints (using the Java program as an example):

Revise

```
public static void SimulateOnce(long seed, long stop)
    throws IOException { .....
```

to

```
public static void SimulateOnce(long seed, long stop, int slower)
    throws IOException { .....
```

where *slower* is s is (s, S) in the inventory system.

- In the main method/function, call the SimulateOnce method/function with stop = 100 and stop = 10000, respectively in two loops whose loop variable changes from slower = 0 to slower = 60 with increment 1.
- Let c_{setup} = \$1,000, c_{hold} = 25, and c_{short} = 700. Compute the dependent cost in an Excel workbook. Graph the cost versus s for the two stop values.

$$C_{dependent} = c_{setup} \overline{u} + c_{hold} \overline{l}^+ + c_{short} \overline{l}^-$$

Submission: both the program and the Excel workbook.

Summary

- Discrete-Event Simulations: random variate vs. trace
- Revisited SSQ
- Revisited SIS
- Variance reduction and robust estimation