A Simple Inventory System

Lawrence M. Leemis and Stephen K. Park, Discrete-Event Simulation: A First Course, Prentice Hall, 2006

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Discrete or Continuous Variables?

- ► Single Server Queue
 - both arrival and service times are continuous variables
- Simple Inventory System
 - ▶ Input variables (e.g., arrival and service times) are inherently discrete

A Simple Inventory System



- Distributes items from current inventory to customers
- Order items from supplier to replenish the inventory
- Customer demand is discrete
 - Customers do not want a portion of an item
- Balancing holding cost, shortage cost, and ordering cost
- Simple
 - One type of item

Different Inventory Policies

- Transaction Reporting Policy
 - Inventory review after each transaction
 - Significant labor may be required
 - ► Less likely to experience shortage
- ► Periodic Inventory Review
 - Inventory review is periodic
 - ▶ Items are ordered, if necessary, only at review times
 - \triangleright Defined by two parameters (s, S)
 - ▶ s: minium inventory level
 - ► S: maximum inventory level
 - s and S are constnat in time, with $0 \le s \le S$
- Assume periodic inventory review
- ▶ Search for (s, S) that minimize cost

Cost and Assumption

- Inventory System Costs
 - Holding cost: for items in inventory
 - ► Shortage cost: for unmet demand
 - Ordering cost: sum of setup and item costs
 - Setup cost: fixed cost when order is placed
 - Item cost: per-item order cost
- Additional Assumptions
 - Back ordering is possible
 - No delivery lag
 - Initial inventory level is S
 - Terminal inventory level is S

Specification Model

- ▶ Time begins at t = 0
- Review times are $t = 0, 1, 2, \dots$
 - ▶ *i*-th time interval begins at time t = i 1 and ends at t = i
- \triangleright l_{i-1} : inventory level at beginning of the *i*-th interval
- ▶ o_{i-1} : amount ordered at time t = i 1 is an integer, $o_{i-1} \ge 0$
- \rightarrow d_i : demand quantity during the i-th interval, $d_i > 0$
- Inventory at end of interval can be negative
 - As a result of back-ordering

Inventory Level Consideration

- ▶ Inventory level is reviewed at t = i 1
- ▶ If $l_{i-1} \ge s$, no order is placed; item If $l_{i-1} < s$, inventory is replenished to S

$$o_{i-1} = \begin{cases} 0 & l_{i-1} \ge s \\ S - l_{i-1} & l_{i-1} < s \end{cases}$$
 (1)

- Items are delivered immediately
- ▶ At end of *i*-th interval, inventory diminished by *d_i*

$$I_i = I_{i-1} + o_{i-1} - d_i (2)$$

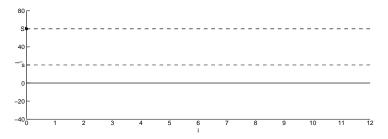
Time Evolution of Inventory Level

Algorithm 1.3.1

```
I_0 = S:
i = 0:
while (more demand to process) {
   i + +:
   if (I_{i-1} < s)
     o_{i-1} = S - I_{i-1}:
   else
     o_{i-1} = 0:
   d_i = GetDemand();
   l_i = l_{i-1} + o_{i-1} - d_i:
n = i;
o_n = S - I_n
I_n = S:
return l_1, l_2, l_3, ... l_n and o_1, o_2, ... o_n;
```

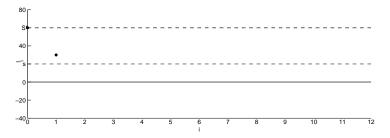
Let (s, S) = (20, 60) and apply Algorithm 1.3.1 to process n = 12 time intervals of operation as follows.

Note that $I_0 = S = 60$.



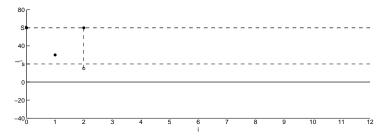
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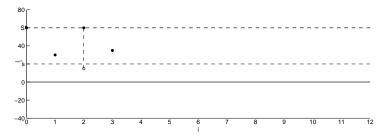
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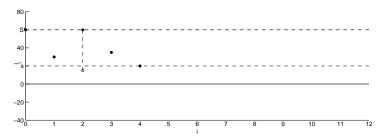
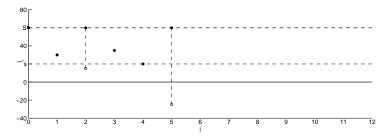


Figure: Inventory Levels

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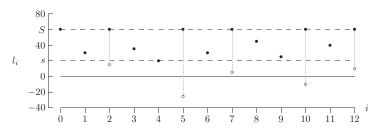


Figure : Inventory Levels

Output Statistics

Average demand and average order per time interval

$$\bar{d} = \frac{1}{n} \sum_{i=1}^{n} d_i \tag{3}$$

$$\bar{o} = \frac{1}{n} \sum_{i=1}^{n} o_i \tag{4}$$

▶ For Example 1.3.1 Data $\bar{d} = \bar{o} = 305/12 \cong 25.42$ items per time interval

In-Class Exercise L3-1

Let (s, S) = (20, 60) and apply Algorithm 1.3.1 by tracing the algorithm to process n = 12 time intervals of operation as follows.

- Calculate average demand. You must show the steps.
- ► Calculate average order. You must show the steps.

Flow Balance

- Average demand and order must be equal
- Ending inventory level is S
- Over the simulated period, all demand is satisfied
- ▶ Average "flow" of items in equals average "flow" of items out



▶ The inventory system is *flow balanced*

Constant Demand Rate Assumption

- Holding and shortage costs are proportional to time-averaged inventory levels
- Must know inventory level for all t
- Assume the demand rate is constant between review times
 - ▶ As a result, the continuous-time evolution of the inventory level is piecewise linear, as illustrated below.

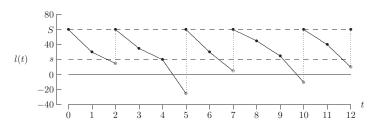


Figure: Piecewise-linear inventory levels

Inventory Level as a Function of Time

▶ Under condition that the demand rate (d_i) is constant between review times, then the inventory level at any time t in i-th interval is,

$$I(t) = I'_{i-1} - (t - i + 1)d_i$$
 (5)

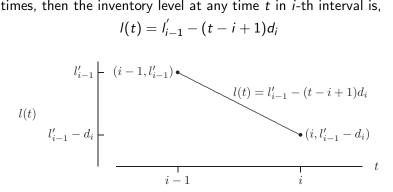


Figure: Linear inventory level in time interval i

 $I'_{i-1} = I_{i-1} + o_{i-1}$ represents inventory level after review.

Inventory Level is Not Linear!

- Inventory level at any time t is an integer
- \triangleright I(t) should be rounded to an integer value
- \blacktriangleright As a result, I(t) is a stair-step, rather than linear, function
- ▶ However, it can be shown, it has no effect on the values of *Time* Averaged Inventory Level \bar{l}_i^+ and \bar{l}_i^-

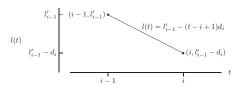


Figure : Linear inventory level in time interval *i*

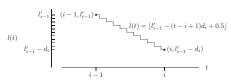


Figure : Linear inventory level in time interval *i*

Time Averaged Inventory Level at Time Interval i

- \triangleright I(t) is the basis for computing the time-averaged inventory level
 - ▶ Case 1: If I(t) remains non-negative over i-th interval

$$\bar{l}_{i}^{+} = \int_{i-1}^{i} l(t)dt \tag{6}$$

▶ Case 2: If I(t) becomes negative at some time τ

$$\bar{I}_{i}^{+} = \int_{i-1}^{\tau} I(t)dt$$
 (7) $\bar{I}_{i}^{-} = -\int_{\tau}^{i} I(t)dt$ (8)

where \bar{l}_i^+ is the time-averaged *holding level* and \bar{l}_i^- is the time-averaged *shortage level*

Case 1: No Back-Ordering

lacktriangle No shortage during i-th time interval if and only if $d_i \leq l_{i-1}^{'}$

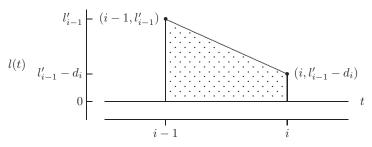


Figure: Inventory level in time interval i without back-ordering

▶ Time-averaged holding level: area of a trapezoid

$$\bar{l}_{i}^{+} = \int_{i-1}^{i} l(t)dt = \frac{l_{i-1}' + (l_{i-1}' - d_{i})}{2} = l_{i-1}' - \frac{1}{2}d_{i}$$
 (9)

Case 2: Back-Ordering

Inventory I(t) becomes negative if and only if $d_i > l'_{i-1}$, i.e., at $t = \tau = i - 1 + l'_{i-1}/d_i$

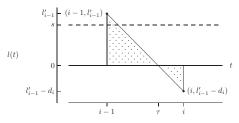


Figure: Inventory level in time interval i with back-ordering

► Time-averaged holding and shortage levels for *i*-th interval computed as the areas of triangles

$$\bar{l}_{i}^{+} = \int_{i-1}^{\tau} l(t)dt = \dots \frac{(l_{i-1}^{\prime})^{2}}{2d_{i}} \quad (10) \quad \bar{l}_{i}^{-} = -\int_{\tau}^{i} l(t)dt = \dots = \frac{(d_{i} - l_{i-1}^{\prime})^{2}}{2d_{i}} \quad (11)$$

\bar{I}_i^+ and \bar{I}_i^-

$$\bar{I}_{i}^{+} = \int_{i-1}^{\tau} I(t)dt = \frac{1}{2}[\tau - (i-1)]I_{i-1}' = \frac{(I_{i-1}')^{2}}{2d_{i}}$$

$$\bar{I}_{i}^{-} = -\int_{\tau}^{i} I(t)dt = -\frac{1}{2}(i-\tau)(I_{i-1}' - d_{i})$$

$$= -\frac{1}{2}[i - (i-1+I_{i-1}'/d_{i})](I_{i-1}' - d_{i})$$

$$= -\frac{1}{2}(1-I_{i-1}'/d_{i})(I_{i-1}' - d_{i})$$

$$= -\frac{1}{2}\frac{d_{i} - I_{i-1}'}{d_{i}}(I_{i-1}' - d_{i})$$

$$= -\frac{(d_{i} - I_{i-1}')(I_{i-1}' - d_{i})}{2d_{i}} = \frac{(d_{i} - I_{i-1}')^{2}}{2d_{i}}$$

Time-Averaged Inventory Level

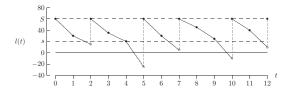


Figure: Piecewise-linear inventory levels

► Time-averaged holding level and time-averaged shortage level

$$\bar{I}^{+} = \frac{1}{n} \sum_{i=1}^{n} \bar{I}_{i}^{+}$$
 (12) $\bar{I}^{-} = \frac{1}{n} \sum_{i=1}^{n} \bar{I}_{i}^{-}$ (13)

- ▶ Note that time-averaged shortage level is positive
- ▶ The time-averaged inventory level is

$$\bar{I} = \frac{1}{n} \int_{0}^{n} I(t) dt = \bar{I}^{+} - \bar{I}^{-}$$
 (14)

In-Class Exercise L3-2

Let (s, S) = (20, 60) and process n = 3 time intervals of operation as follows (see slides 23, 24, and 26 for steps)

$$i$$
 | 1 | 2 | 3 | 15 | 45 | 30 |

- Calculate time-averaged holding level
- ► Calculate time-averaged shortage level

Computational Model

- Program sis1 is a trace-driven computational model of the SIS
- Computes the statistics
 - ightharpoonup d: Average demand
 - ▶ ō: Average order
 - $ightharpoonup \overline{I}^+$: Time-averaged holding level
 - $ightharpoonup \overline{I}^-$: Time-averaged shortage level

and the order frequency \bar{u}

$$\bar{u} = \frac{number\ of\ orders}{n}$$

▶ Consistency check: compute \bar{d} and \bar{o} separately, then compare. They should be equal.

Example 1.3.4: Executing sis1

- ▶ Trace file *sis1.dat* contains data from n = 100 time intervals
- ▶ Inventory-policy parameter values (i.e., mininum & maximum inventory levels) (s, S) = (20, 80)
- Program outputs:

$$\bar{o} = \bar{d} = 29.29$$
 $\bar{u} = 0.39$ $\bar{l}^+ = 42.90$ $\bar{l}^- = 0.25$

In-Class Exercise L3-3: sis1

- ► Run either C/C++ or Java program against the trace with the following inventory levels,
 - $S = 80 \text{ and } s = 2, 4, 6, \dots 40$

submit the result.

Operating Cost

- A facility's cost of operation is determined by
 - c_{item}: unit cost of new item
 - c_{setup}: fixed cost for placing an order
 - chold: cost to hold one item for one time interval
 - c_{short}: cost of being short one time for one time interval

Example 1.3.5: Case Study

- ▶ An automobile dealership that uses weekly periodic inventory review
 - ► The facility is the showroom and surrounding areas holding cars
 - ▶ The items are cars for sell
 - ▶ The supplier is the car manufacturer
 - ▶ The customers are the people who purchase cars from the dealership
 - Assume SIS: sells one type of car

Example 1.3.5: Results of Case Study

- Limited to a maximum of S = 80 cars
- \blacktriangleright Limited to a minimum of s=20 cars
- Inventory reviewed every Monday
 - if inventory falls below s, order cars sufficient to restore the inventory to S
- ► For now, ignore delivery lag
- Then costs:
 - Item cost is $c_{item} = \$8,000$ per item
 - Setup cost is $c_{setup} = \$1,000$ per order from manufacturer
 - ▶ Holding cost is $c_{hold} = 25 per week
 - ▶ Shortage cost is $c_{hold} = 700 per week

Per-Interval Average Operating Costs

- ▶ The average operating costs per time interval are
 - $ightharpoonup c_{item} \bar{o}$: item cost
 - $ightharpoonup c_{setup} \bar{u}$: setup cost
 - $\sim c_{hold}I_{-}^{+}$: holding cost
 - ▶ c_{short} I⁻: shortage cost
- ▶ The average total operating cost per time interval is their sum

$$\bar{C}_{total} = c_{item}\bar{o} + c_{setup}\bar{u} + c_{hold}\bar{l}^{+} + c_{short}\bar{l}^{-}$$
 (15)

► Total cost of operation is the product of the average total operating cost per time interval and the number of intervals

$$C_{total} = n\bar{C}_{total} \tag{16}$$

Example 1.3.6: Operating Cost

▶ Use the statistics in Example 1.3.4

$$\bar{o} = \bar{d} = 29.29$$
 $\bar{u} = 0.39$ $\bar{l}^+ = 42.90$ $\bar{l}^- = 0.25$

and the constants in Example 1.3.5

$$c_{item} = \$8,000$$
 $c_{setup} = \$1,000$ $c_{hold} = \$25$ $c_{hold} = \$700$

- For the dealership
 - item cost: $\$8,000 \times 29.29 = \$234,320$ per week
 - setup cost: $\$1,000 \times 0.39 = \390 per week
 - ▶ holding cost: $$25 \times 42.40 = $1,060$ per week
 - ▶ shortage cost: $$700 \times 0.25 = 175 per week

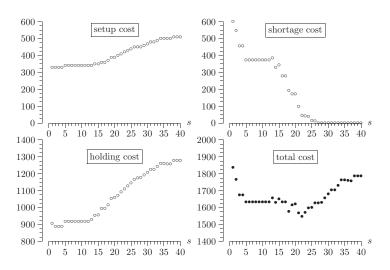
Cost Minimization

- ▶ By varying minimum inventory level s (and possibly maximum inventory level S), an optimal policy can be determined
- ▶ Optimal ←⇒ minimum average total cost per time interval
- Note that $\bar{o} = \bar{d}$ and \bar{d} depends only on the demands
- ▶ Hence, item cost per time interval $c_{item}\bar{o}$ is independent of (s, S)
- Average dependent cost per time interval is the sum of these three
 - $ightharpoonup c_{setup} \bar{u}$: average setup cost per time interval
 - $ightharpoonup c_{hold}\bar{l}^+$: average holding cost per time interval
 - $c_{short}\bar{l}^-$: average shortage cost per time interval

Experiment

- ▶ Let S be fixed, and let the demand sequence be fixed
- ▶ If *s* is systematically increased, we expect:
 - Average setup cost and holding cost per time interval will increase as s increases
 - ▶ Average shortage cost per time interval will decrease as s increases
 - Average dependent cost per time interval will have "U" shape, yielding an optimum (i.e., minimum cost = \$1,550 at s = 22)

Example 1.3.7: Optimal Periodic Inventory Review Policy



• Minimum cost = \$1,550 at s = 22

In-Class Exercise L3-4: sis1 and Optimizing Operating Cost

- ▶ Use the sample data (sis1.dat) and the parameters in the lecture nodes, verify the optimal inventory review policy by showing the results in slides 35 37, produce the data necessary to reproduce the figures in slide 38.
- Submit the result

Summary

- ► SIS
- Cost model and case studies