

# The Internet Protocol

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# Outline

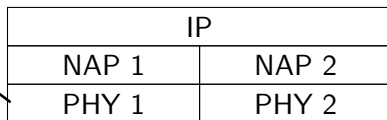
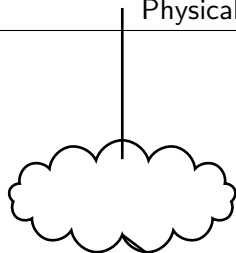
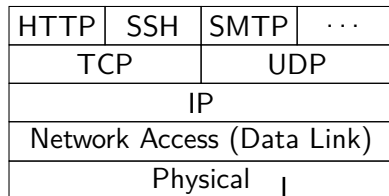
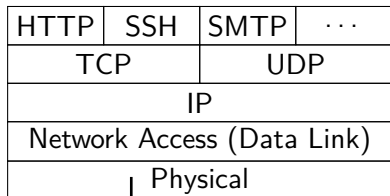
- 1 Overview of Internetworks
- 2 Overview of Internet Protocol
- 3 IPv4
  - IPv4 Address
  - IPv4 Packet Forwarding
  - IPv4 Packet Format
  - Control and Error Reporting
  - IP and Network Access Protocol
  - Host Configuration
- 4 IPv6

# The Internet and Internetworking

The Internet is a network of networks, i.e., an internetwork.

- ▶ What are the motivation to build such a network?
- ▶ What are the requirements?
- ▶ What are the problems?

# TCP/IP Internetworks



# Challenges

- ▶ (Heterogeneity) The networks in an internetwork are heterogeneous.
- ▶ (Scalability) There are a lot of hosts and networks on an internetwork.
- ▶ (Survivability) Hosts on part of networks should still be able to communicate when others fail.
- ▶ ...

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# Internet Protocol Operation

The Internet Protocol (IP) provides a connectionless service (or datagram service) between two end systems (e.g., two hosts, or two routers, or a router and a host)

- ▶ It is easy to deal with variety of networks as it does not require much for these networks, and it is simple by itself.
- ▶ The connectionless service delivers packets with the best effort approach and there is no centralized control, which makes IP service is highly robust.
- ▶ It is best for connectionless transport protocols, as it does not impose unnecessary overhead.
- ▶ IP realizes a network of networks, which allows a network (i.e., an internetwork) to scale.

# Design Issues

- ▶ Routing
- ▶ Datagram lifetime
- ▶ Fragmentation and assembly
- ▶ Error control
- ▶ Flow control
- ▶ Congestion control



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# IPv4 Address

An IPv4 address is 4 bytes (or 32 bits).

- ▶ Unicast address ( $00000000_h/1$  ,  $80000000_h/2$ , and  $C0000000/3$ )
- ▶ Multicast address ( $E0000000_h/4$ )
- ▶ Reserved address ( $F0000000_h/5$ )
- ▶ Broadcast address

# IPv4 Unicast Address

- ▶ Private network
  - ▶ 192.168.0.0/16
  - ▶ 172.16.0.0/12
  - ▶ 10.0.0.0/8
- ▶ Loopback. 127.0.0.0/8
- ▶ Link-local unicast. 169.254.0.0/16
- ▶ Documentation (TEST-NET-1). 192.0.2.0/24
- ▶ Documentation (TEST-NET-2). 198.51.100.0/24
- ▶ Documentation (TEST-NET-3). 203.0.113.0/24
- ▶ Global unicast. Everything else with exceptions.

# IPv4 Unicast Address: Network and Host Parts

An IPv4 address consists of two parts:

- ▶ Network part
- ▶ Host part

To understand this, let's look at the IP forwarding algorithm

# IPv4 Packet Forwarding

Conceptually, we can describe the IPv4 packet forwarding (unicast packet forwarding) algorithm as follows,

```
D = destination IP address // D = <NetworkNumber, HostNumber>
Selected = {}
for each entry of <NetworkNumber, NextHop> in forwarding table
    // Get_Network_Number returns the entry
    // if D.NetworkNumber = NetworkNumber
    Ds = Get_Network_Number(D) // {Ds} can be  $\emptyset$  if not a match
    Selected = Selected  $\cup$  {Ds}
endfor
Ds = Longest_Match(Selected)
if {Ds}  $\neq \emptyset$ 
    if Ds.NextHop is an interface
        deliver datagram directly to destination
    else
        deliver datagram to Ds.NextHop (a router)
else
    drop the packet
```

# Forwarding Table

- ▶ A forwarding table consists of entries of network and next hop entity where a next hop entity is either a network interface or a next hop router.
- ▶ Run these to observe examples.

```
ip route show
```

Listing: Linux

```
netstat -n -r
```

Listing: Unix and Windows

# Fragmentation and Reassembly

- ▶ Different network can have different MTU
- ▶ MTU. Maximum transmission unit is the maximum sized datagram that can be transmitted through the next network is called the maximum transmission unit.
  - ▶ IPv4 over Ethernet V2. MTU = 1500 bytes.
  - ▶ IPv4 over IEEE 802.11 Wi-Fi (WLAN). MTU = 2304 bytes before encryption.
  - ▶ IPv4 over FDDI. MTU = 4352 bytes.
- ▶ But how big an IPv4 packet can be?
- ▶ Fragmentation (<https://tools.ietf.org/html/rfc791#page-26>)
- ▶ Reassembly (<https://tools.ietf.org/html/rfc791#page-28>)

# IPv4 Packet Header

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
V=4				IHL				DS				ECN <sup>1</sup>		Total Length																	
Identification														Flags			Fragment Offset														
Time To Live						Protocol <sup>2</sup>						Header Checksum																			
Source Address																															
Destination Address																															
Options + Paddings (32 bits × n, n = 0, 1, 2, ...) ...																															

<sup>1</sup>See RFC 3168

<sup>2</sup><https://www.iana.org/assignments/protocol-numbers>



# Internet Control and Message Protocol (ICMP)

- ▶ ICMP is a network layer protocol, is also a user of IP, i.e., an ICMP packet is the pay load of an IP packet. (RFCs 792, 950, 1250, 1393, 1475, 1788)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
Type								Code								Checksum															
Varying based on Type																															
...																															

## Example ICMP Message Types

- ▶ Type 0. Echo Reply Message
- ▶ Type 8. Echo (Request) Message
- ▶ Type 3. Destination Unreachable Message
- ▶ Type 5. Redirect Message
- ▶ Type 30. Traceroute Message (RFC 1393)
- ▶ Type 11. Time Exceeded

# Address Resolution

- ▶ Needs to establish the mapping between IP address of the next-hop node and the hardware (or MAC) address
- ▶ Address Resolution Protocol (ARP).
- ▶ Let's try,

```
sudo arp
```

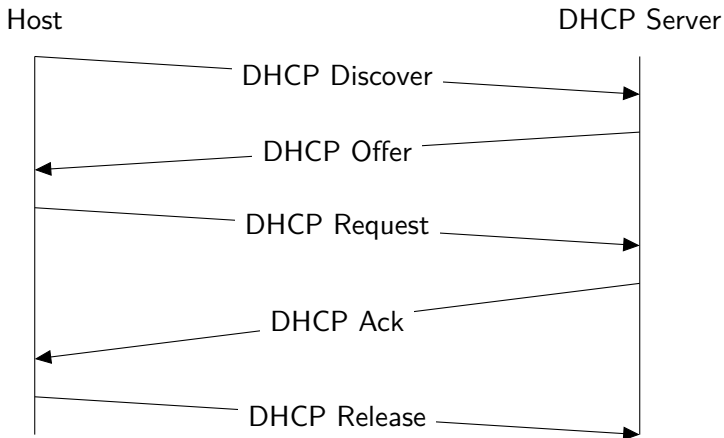
Listing: Linux

- ▶ Where do these come from?

# Dynamic Host Configuration Protocol (DHCP)

- ▶ How do we assign IPv4 addresses?
- ▶ We can enable dynamic allocation of IPv4 addresses and supply other configuration parameters to hosts (automatically) via DHCP (RFC 2131).
- ▶ DHCP is a user of UDP, i.e., DHCP uses UDP as its transport protocol.

## Example of DHCP Message Exchange



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# Comparing with IPv4

- ▶ IPv6 address
- ▶ Packet format
- ▶ Address translation and host configuration?

# IP Next Generation to IPv6

To apply a few important findings from IPv4 experience and research,

- ▶ IPv4 Address space (32 bits) is too small. IPv6 address space is 128 bits.
- ▶ Routers often do not handle optional fields in IPv4 headers. IPv6 options (extension headers) are placed in separated optional headers, which simplifies processing delay at routers
- ▶ Do away with DHCP? IPv6 has address autoconfiguration capability and can assign addresses without an application layer server like DHCP.
- ▶ Support a few innovations? IPv6 introduces anycast address to support anycast, and adds a scope field to multicast address to improve scalability of multicast routing.
- ▶ Resource allocation is difficult in IPv4. IPv6 can label packets for different traffic flow, which aids resource allocation for specialized traffic (e.g., real-time video)

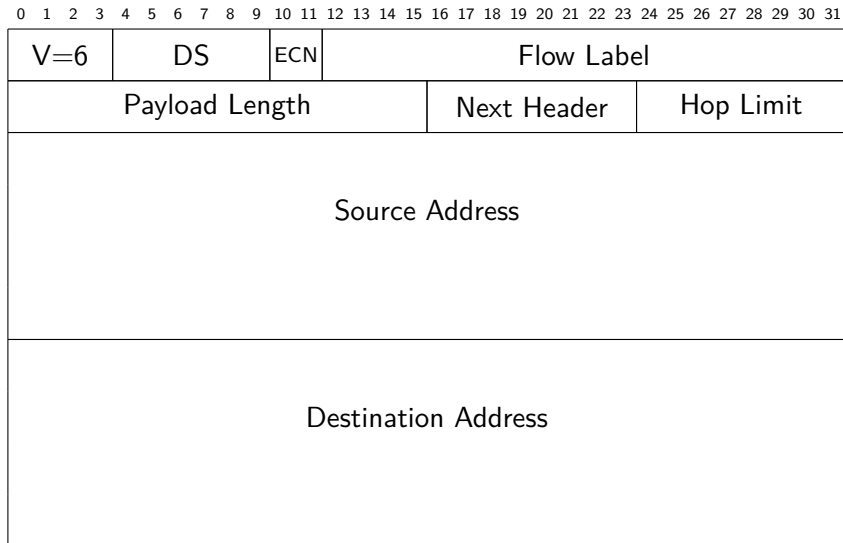


# IPv6 Address

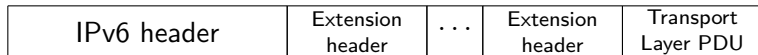
Divided into a few categories, e.g.,

- ▶ Embedded IPv4 Address ( $:: FFFF/96$ )
- ▶ Loopback ( $:: 1/128$ )
- ▶ Link-local unicast ( $FE80 :: /10$ )
- ▶ Private network. ( $FC00 :: /7$ )
- ▶ Documentation ( $2001 : 0DB8 :: /32$ )
- ▶ Multicast ( $FF00 :: /8$ )
- ▶ Global unicast (Everything else with exceptions)

# IPv6 Packet Header



# IPv6 Packet Structure



# IPv6 Stateless Address Autoconfiguration

- ▶ Assign IPv6 addresses to hosts without a server? (RFC 4862)
  1. Generate a link-local address.
  2. Test the uniqueness of a link-local address.
  3. Assign a link-local address.
  4. Contact the router.
  5. Provide direction to the node.
  6. Configure the global address.