

Shortest Path Algorithms

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Dijkstra's Algorithm

Algorithm 1: Dijkstra's Shortest Path Algorithm

Input : G , Graph $G = (V_G, E_G)$

Input : $s, s \in V_G$ is the source node of the shortest paths in Graph G

Function DijkstraShortestPath(G, s)

```
for  $n \in V_G$  do
     $d(s, n) \leftarrow \infty$ 
     $p(n) \leftarrow nil$ 
     $N_G = V_G$ 
 $d(s, s) \leftarrow 0$ 
while  $N_G \neq \phi$  do
     $u \leftarrow u | d(s, u) = \min_{\forall n \in N_G} d(s, n)$ 
     $N_G \leftarrow N_G \setminus \{u\}$ 
    for  $v \in \text{Neighbor}(G, u)$  do
         $l \leftarrow d(s, u) + w(u, v)$ 
        if  $l < d(s, v)$  then
             $d(s, v) \leftarrow l$ 
             $p(v) \leftarrow u$ 
return  $d, p$ 
```

Bellman-Ford Algorithm

Algorithm 2: Bellman-Ford Shortest Path Algorithm

Input : G , Graph $G = (V_G, E_G, W_G)$

Input : s , $s \in V_G$ is the source node of the shortest paths in Graph G

Function BellmanFordShortestPath(G, s)

```
for  $n \in V_G$  do
   $d(s, n) \leftarrow \infty$ 
   $p(n) \leftarrow nil$ 
 $d(s, s) \leftarrow 0$ 
for  $u \in V_G$  do
  for  $(u, v) \in E_G$  do
     $l \leftarrow d(s, u) + w(u, v)$ 
    if  $l < d(s, v)$  then
       $d(s, v) \leftarrow l$ 
       $p(v) \leftarrow u$ 

for  $(u, v) \in E_G$  do
  if  $d(s, v) > d(s, u) + w(u, v)$  then
    RaiseError("Negative Cycle")

return  $d, p$ 
```