

CISC 7332X T6

# Collision-free and Limited-contention Protocols

Hui Chen

Department of Computer & Information Science

CUNY Brooklyn College

# Outline

- Channel allocation problem
- Multiple Access Protocols
  - Discussed
    - ALOHA
    - CSMA (Carrier Sense Multiple Access)
    - CSMA/CD
  - Collision-free protocols
  - Limited-contention protocols
  - To discuss next week
    - Ethernet protocols
    - Wireless LAN protocols

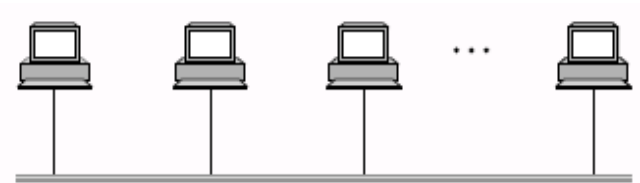
# Medium Access Control

- Two types of network links

- Point-to-point



- Multiple access (broadcast)

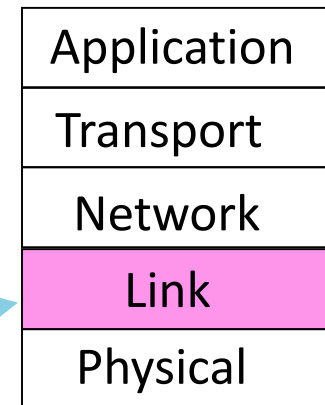


- Key issue

- Who gets to use the channel when there is a competition to it?
- Multiaccess channel/random access channel
- Medium Access Control (MAC)

# The MAC Sublayer

- The protocols used to determine who goes next on a multiaccess channel
- Especially important for LAN, particularly wireless LANs
- In contrast, WANs general use point-to-point links, excepts for satellite networks



MAC is in here!

# Multiple Access Protocols

- Discussed
  - ALOHA
  - CSMA (Carrier Sense Multiple Access)
- Collision-free protocols
- Limited-contention protocols
- To discuss next week
  - Ethernet protocols
  - Wireless LAN protocols

# Collision Free Protocols

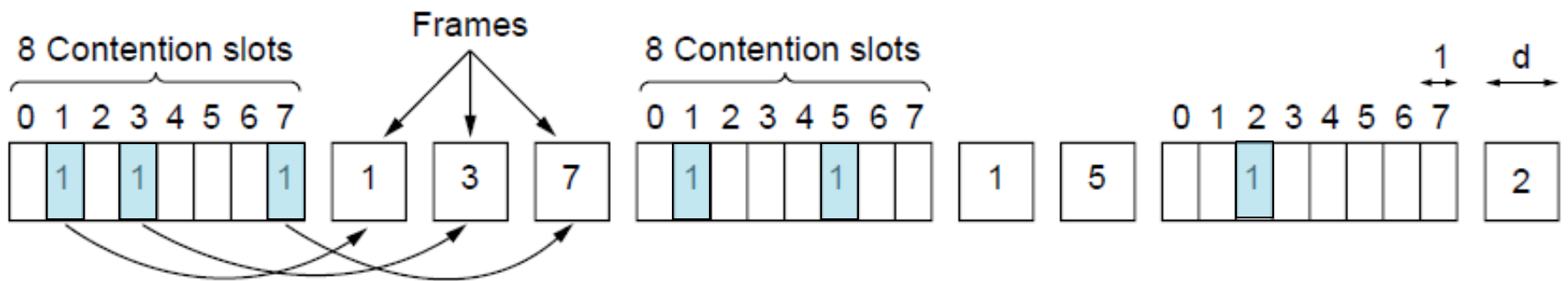
- Collision-free protocols avoid collisions entirely
  - Senders must know when it is their turn to send
- Bitmap

# Bitmap

- The basic bit-map protocol:
  - The protocol consists of two phases
    - Contention
      - There are a fixed number of content slots.
      - Sender set a bit in contention slot if they have data
    - Sending frames
      - Senders send in turn; everyone knows who has data
- Protocols like this, in which, the desire to transmit is broadcast before actual transmission is called reservation protocols

# Basic Bitmap: Example

- Each station announces whether it has data frame to transmit in *its* content slot
- After contention slots, transmit if any station has a frame to transmit





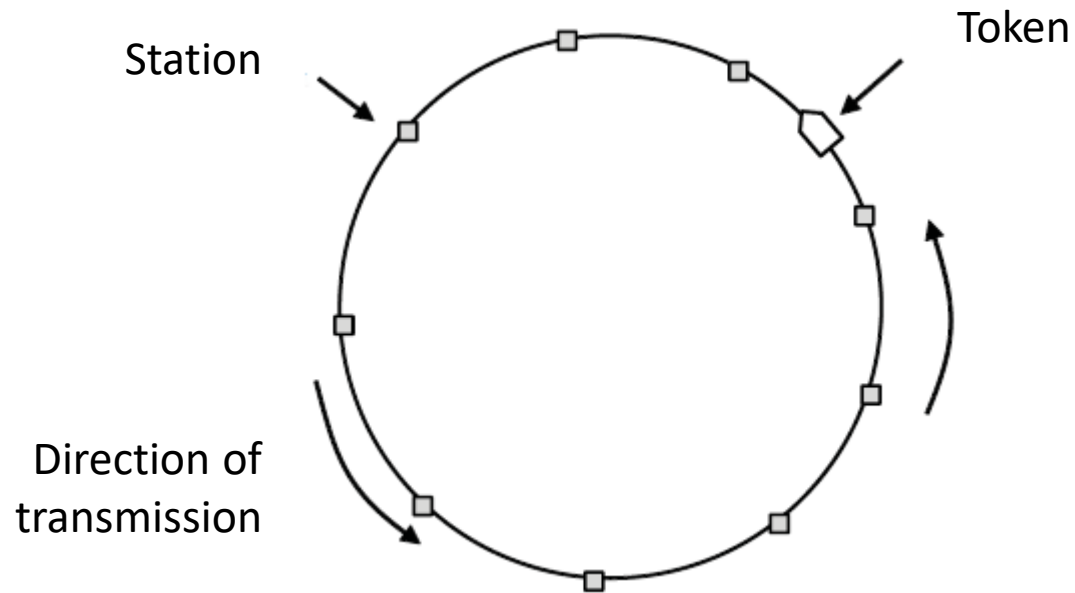
# Basic Bitmap: Efficiency

- Channel efficiency/utilization: assume  $N$  slots,  $d$  bits per frame
  - Low load:  $\sim d/(d+N)$
  - High load:  $\sim d/(d + 1)$
- Mean delay
  - $> (N-1)d/2 + N$

# Token Passing

- Token sent round ring defines the sending order
  - Station with token may send a frame before passing
  - Idea can be used without ring too, e.g., token bus

# Token Passing: Example



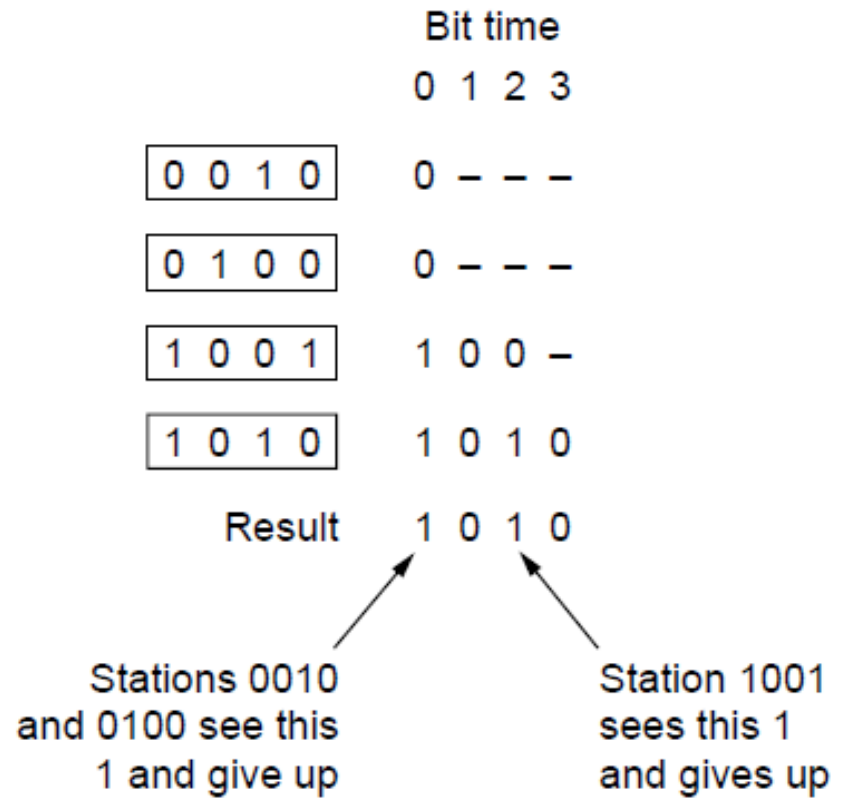
# Token Passing: Example Protocols

- IEEE 802.5/Token Ring
- FDDI (Fiber Distributed Data Interface)
- IEEE 802.17/RPR (Resilient Packet Ring)

# Count Down

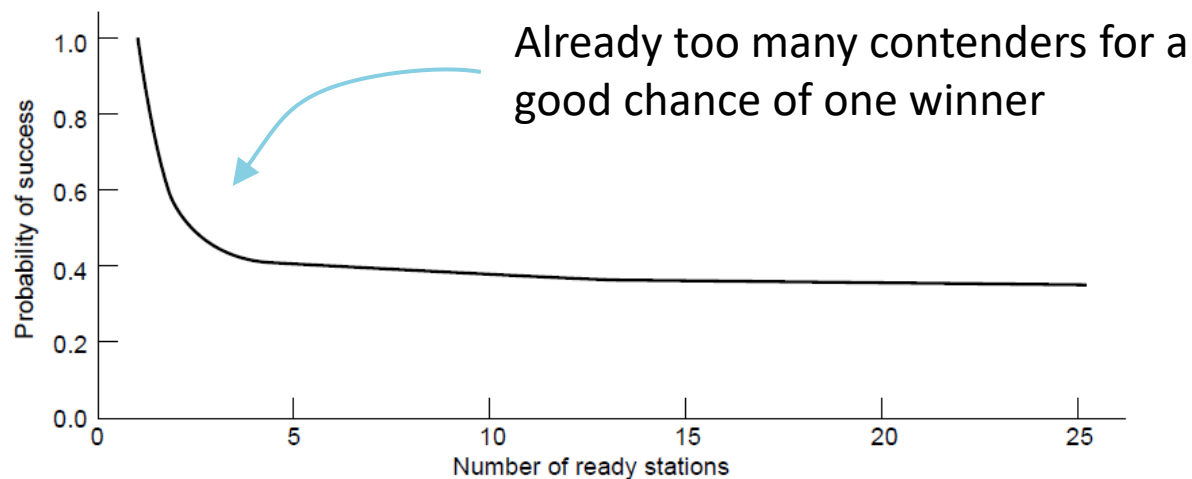
- An improvement over basic Bitmap
- Stations send their address in contention slot (log N bits instead of N bits), one bit at a time
- Medium ORs bits; stations give up when they send a "0" but see a "1"
- Station that sees its full address is next to send
- Requirement on physical layer?
  - Transmitting while receiving ...

# Count Down: Example



# Limited-Contention Protocols

- To divide stations into groups within which only a very small number are likely to want to send
  - Avoids wastage due to idle periods and collisions



# Adaptive Tree Walk: Motivating Example

- U.S. Army adopted a method to test soldiers for syphilis during World War II aiming at reducing tests.
  1. Took a blood sample from  $N$  soldiers
  2. Pour a portion of each to a test tube
  3. If the mixed sample is clean with antibodies, the  $N$  soldiers are declared healthy
  4. Otherwise, divided  $N$  into two groups, each with  $N/2$  soldiers
  5. Repeat the testing protocol starting at step 1 (albeit for each group of  $N/2$  soldiers) recursively until the infected soldier is found.

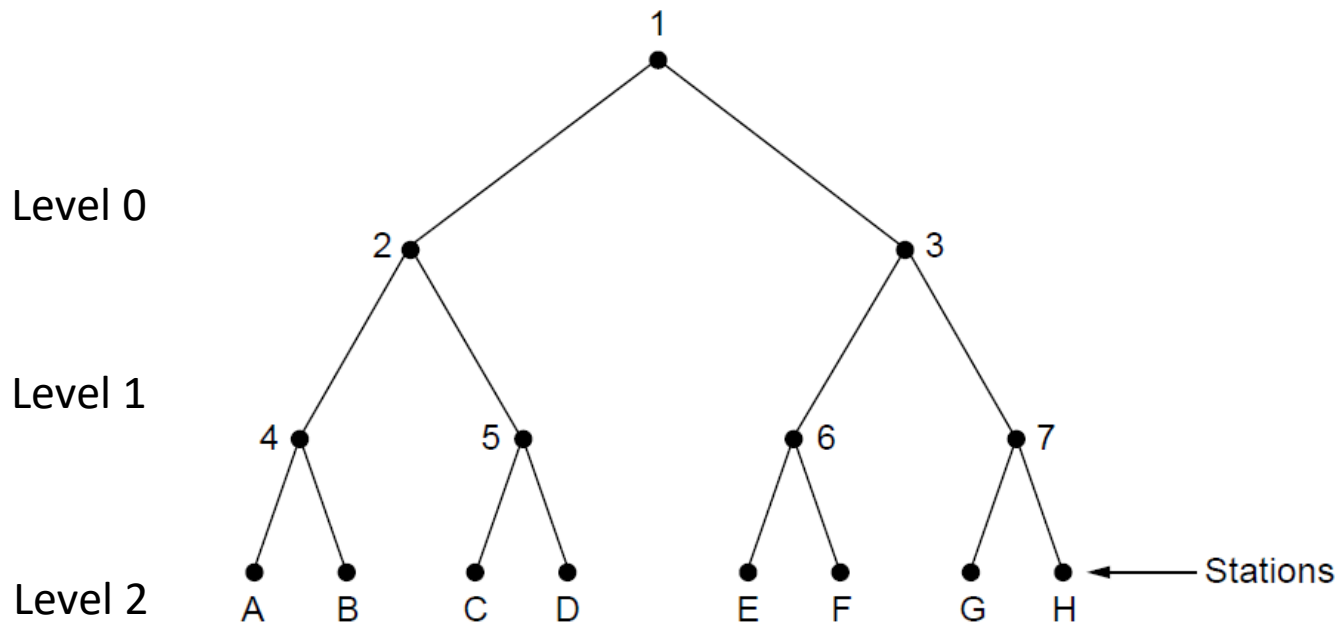


# Limited Contention: Adaptive Tree Walk

- Tree divides stations into groups (nodes) to poll
  - Depth first search under nodes with poll collisions
  - Start search at lower levels if  $>1$  station expected

# Limited Contention: Adaptive Tree Walk: Example Tree

- 8 stations; time slots conveniently numbered 1, 2, ... 7 as the node numbers in the tree; nodes are grouped accordingly
- If contention fails at time slot 1, nodes under 2 contends at slot 2, and nodes under 3 contends at slot 3 ...



# Questions?

- Collision-free protocols
- Limited-contention protocols