

# Process Management

Hui Chen <sup>a</sup>

<sup>a</sup>CUNY Brooklyn College

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# Outline

- 1 Concept of Process and Motivation
- 2 Motivation for *Process*
- 3 Policy and Mechanism
- 4 Efficient Use of Resources
- 5 Designing Process: Main Memory
- 6 Designing Processes: Execution
- 7 Process States Transition and Queues
- 8 OS Design Objectives and CPU Scheduler
- 9 Process Operations and Example Programs
- 10 Querying Process Status on Linux Systems

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# Overview of OS Services

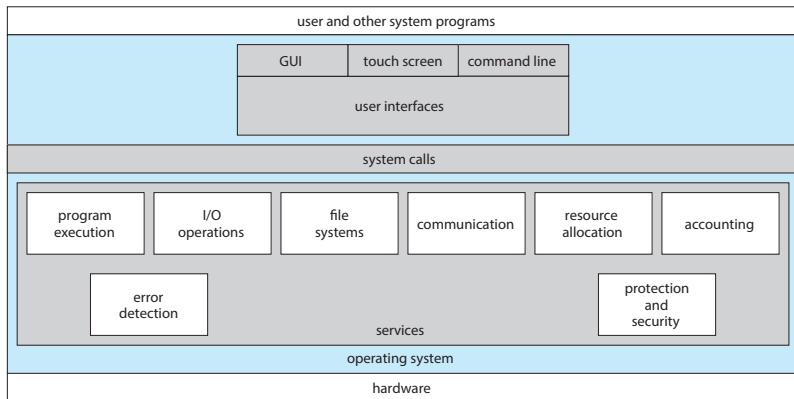


Figure: A view of operating system services<sup>1</sup>.

<sup>1</sup>Abraham Silberschatz, Peter B. Galvin, and Greg Gagne. *Operating system concepts*. 10th edition. John Wiley & Sons, 2018.

## Evolved to *Process*

- ▶ Early computers run a single program at a time.
- ▶ Contemporary computer systems allow multiple programs to be loaded into memory and executed concurrently.
- ▶ This evolution results in the notation of a process, i.e., a program in execution.

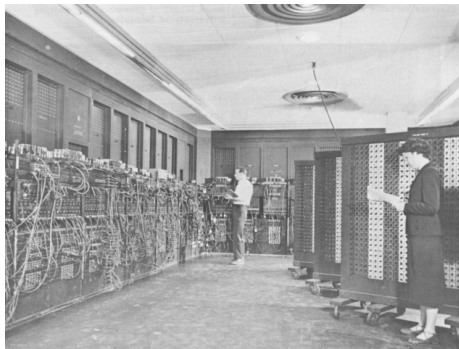


Figure: From [US Army Photo Archive](#)

# Process Management

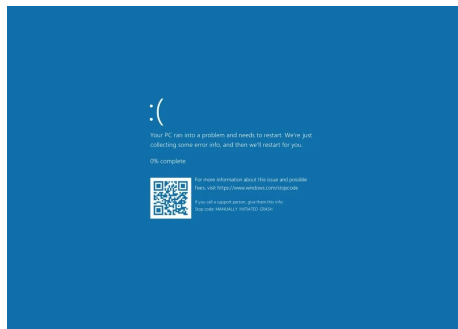
- ▶ Why?
- ▶ What?
- ▶ How?

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# Need for Reliability and Protection

- ▶ Programs can fail.
- ▶ Programs may access other programs' code and data, intentionally or unintentionally.
- ▶ ...





# Need for Efficient Use of Resources

- ▶ Computer systems has rather limited and also “expensive” resources
  - ▶ CPU cycles
  - ▶ Main memory (RAM)
  - ▶ I/O devices (secondary storage, network, ...)
  - ▶ ...

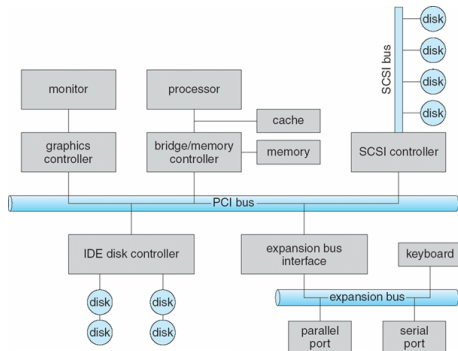


Figure: Source: Figure 12-1 in Silberschatz et al., 2018<sup>a</sup>

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# Policy

- ▶ Policy (What do we want?). The OS should provide *efficient and protected* access to shared resources.
  - ▶ Programs can run concurrently to utilize resources.
  - ▶ No program can read or write memory of another program or of the OS (what memory?).
  - ▶ *Additionally*, do we wish to have better utilization of system resources, e.g., do we wish to have better (small) response time?

# Mechanism

- ▶ Policy (What?). The OS should provide *efficient and protected* access to shared resources.
- ▶ Mechanism (How?). Hardware, OS kernel, and *process*.
  - ▶ Realizing “Virtual” resources, e.g., address translation. Each process has its isolated logical (or virtual address space), and hardware translates virtual address to physical address.
  - ▶ Supporting dual mode operating. Processes run in user or kernel mode, i.e., to restrict process in user mode from accessing privileged instructions (which needs both hardware and software support)

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# Concepts of Multiprogramming and Time-Sharing

- ▶ Multiprogramming
- ▶ Time-sharing

# Multiprogramming

- ▶ A technique in OS that OS organizes a collection of processes in such a way that the processes run concurrently (in parallel or in pseudo-parallel) and the CPU always has a process to execute.
- ▶ Multiprogramming can increase CPU utilization.
- ▶ Why?

# Time-Sharing

- ▶ A technique in OS that OS used a timer and cycle processes rapidly through the CPU, giving each user a share of the resources.
- ▶ Time-sharing can decrease response time to make system more responsive.



# A Simple Model of Multiprogramming

$$\text{CPU Utilization} = 1 - p^n \quad (1)$$

where

- ▶  $n$ . There are  $n$  processes in memory, and we call it the degree of multiprogramming.
- ▶  $p$ . The fraction of time that a process spends in waiting for I/O to complete.

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# Process in Memory

OS creates an illusion that a process runs as if the process ran on its own CPU and had its own memory, i.e., a process gets its own

- ▶ “virtual CPU”, and
- ▶ “virtual memory address space”

## Layout of Process in Memory

The layout of a process in memory consists of multiple parts, generally,

- ▶ Text. The text section is the program code.
- ▶ Stack. The stack section contains temporary data, such as, function parameters, return addresses, local variables.
- ▶ Data. The data section containing global variables
- ▶ Heap. The heap section containing memory dynamically allocated during run time.

# Memory Layout of a C Program

Let's take a look at a C program ...

## Multiple Processes in Memory

Consider that there are now multiple processes in memory.

- ▶ But there is only one single physical memory address space in the machine, how do we create the *illusion* that each has its own memory address space?

# Address Translation

A process operates in the memory address space distinct from the physical memory space of the machine

- ▶ Address translation. Hardware translates a memory address the process operates on to a memory address in the physical memory.
- ▶ Shall discuss in detail later ...

## Recall Discussion: Kernel and Process

- ▶ Policy. The OS should provide efficient and protected access to shared resources.
  - ▶ No program can read or write memory of another program or of the OS.
  - ▶ Programs can run concurrently to utilize resources efficiently.
- ▶ Mechanism. Hardware, OS kernel, and process.
  - ▶ Dual mode operating. Processes run in user or kernel mode. Restrict process in user mode from accessing privileged instructions.
  - ▶ Address translation. Each process has its isolated logical (or virtual address space), and hardware translates virtual address to physical address.



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## Running Multiple Processes?

Consider that two processes run concurrently (in pseudo-parallel) in a single CPU core system ...

- ▶ These two processes switch on and off the CPU for many rounds ...  
how do we create the *illusion* that each has its own CPU?

## Two processes are running ...

Consider

- ▶ At  $T_1$ , Process  $P_1$  is active (on the physical CPU), and Process  $P_2$  is inactive.
- ▶ At  $T_2$ , Process  $P_2$  is active (on the physical CPU), and Process  $P_1$  is inactive.

## Two processes are running ...

How?

1. OS maintains a data structure in memory (its own address space), when the OS runs, it switches  $P_1$  off the CPU and  $P_2$  on the CPU, i.e., it
2. saves registers Program Counter (PC) and Stack Pointer (SP) in the data structure for  $P_1$  in memory
3. loads PC and SP from the data structure for  $P_2$  in memory

## Supporting Data Structure: Process Control Block

A process control block (also called task control block or thread control block) is a data structure for information associated with each process including the process's execution context.

- ▶ Process state, i.e., running, waiting, etc
- ▶ Program counter, i.e., location of instruction to next execute
- ▶ Stack pointer, i.e., top of the process stack
- ▶ other CPU registers if necessary, i.e., contents of all process-centric registers
- ▶ CPU scheduling information, i.e., priorities, scheduling queue pointers
- ▶ Memory-management information, i.e., memory allocated to the process
- ▶ Accounting information, i.e., CPU used, clock time elapsed since start, time limits
- ▶ I/O status information, i.e., I/O devices allocated to process, list of open files

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# Queues and Transition

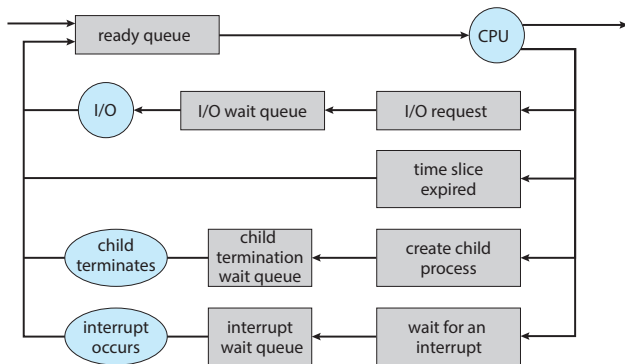


Figure: Process queues and transitions<sup>2</sup>.

<sup>2</sup>Abraham Silberschatz, Peter B. Galvin, and Greg Gagne. *Operating system concepts*. 10th edition. John Wiley & Sons, 2018.

# Data Structures for Process Queues

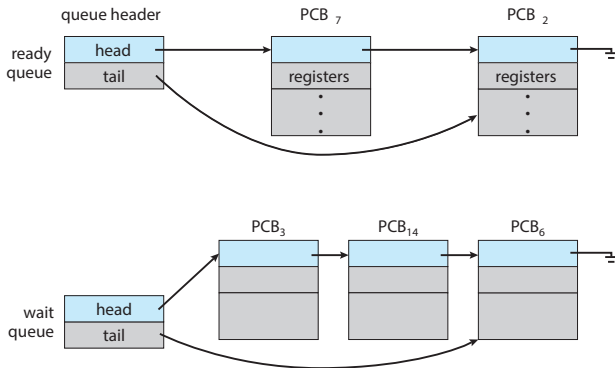


Figure: The *ready* and *wait* queues<sup>3</sup>.

<sup>3</sup>Abraham Silberschatz, Peter B. Galvin, and Greg Gagne. *Operating system concepts*. 10th edition. John Wiley & Sons, 2018.



# Process States and Transition

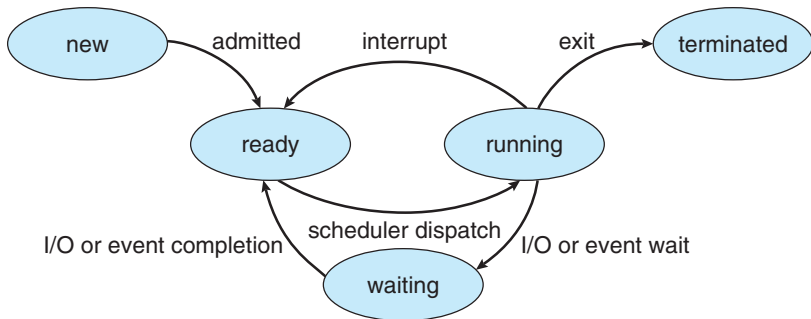


Figure: Process transitions in an OS<sup>4</sup>.

<sup>4</sup>Abraham Silberschatz, Peter B. Galvin, and Greg Gagne. *Operating system concepts*. 10th edition. John Wiley & Sons, 2018.

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# Multiprogramming and Time-Sharing

- ▶ The objective of multiprogramming is to have some process running at all times to maximize CPU utilization.
- ▶ The objective of time-sharing is to switch a CPU core among processes so frequently that users can interact with each program while it is running.

Questions. How do we support these objectives?

# CPU Scheduler

- ▶ The CPU scheduler is to select from among the processes that are in the ready queue and allocate a CPU core to one of them.
- ▶ Different CPU scheduling algorithms help OSES meet different objectives.

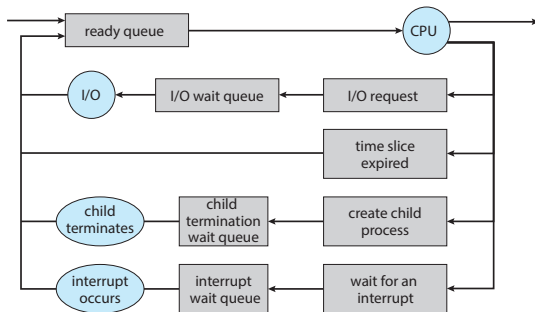


Figure: Process queues and transitions<sup>5</sup>.

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# Process Operations

- ▶ An OS must provide services for,
  - ▶ process creation,
  - ▶ process termination, and
  - ▶ others
- ▶ Design considerations and related concepts.

# System Calls for Process Operations in UNIX

- ▶ Process creation. `fork()`, `exec()`, `wait()` ...
- ▶ Process termination. `exit()`, `wait()`, `abort()` ...

# Multiprocess Architecture

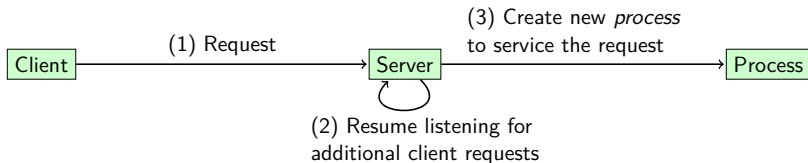


Figure: A multiprocess architecture server program



# Design Multiprocess Architecture Programs

Design programs of multiprocess architecture on UNIX ...

- ▶ Observe the example programs and consider rationale behind the design ...

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# Examining Process States on Linux Systems

- ▶ Recall previous discussion ...
  - ▶ `ps`, `pstree`, `top`, `vmstat`, `iostat`, `lsof`  
`,` ...
- ▶ The `/proc` file system  
`man proc`

# References I



Silberschatz, Abraham, Peter B. Galvin, and Greg Gagne. *Operating system concepts*. 10th edition. John Wiley & Sons, 2018.