

# CPU Scheduling

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# Outline

- 1 CPU Scheduling
- 2 Scheduling Criteria
- 3 Scheduling Algorithms
- 4 Thread Scheduling and Multiprocessor Scheduling

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# Recall Process Queues and State Transitions

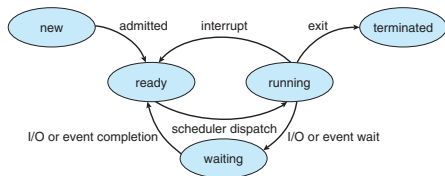
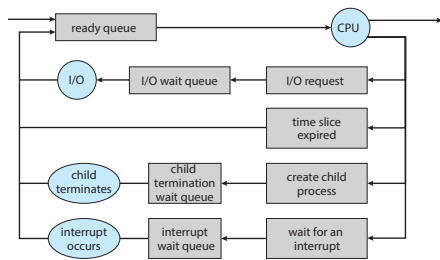


Figure: Process queues and transitions<sup>1</sup>.

- ▶ CPU scheduling is the basis of multiprogrammed operating systems, it is about selecting a task from the *Ready* queue to execute it on CPU.
  - ▶ Process scheduling vs. thread scheduling

<sup>1</sup>Silberschatz, Galvin, and Gagne, *Operating system concepts*.

# When does scheduling happens?

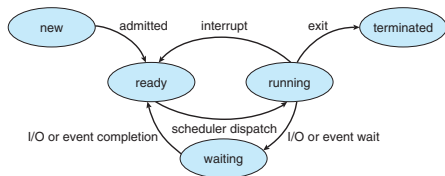
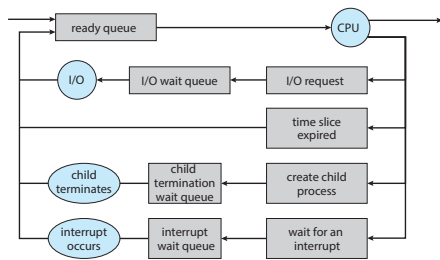


Figure: Process queues and transitions<sup>2</sup>.

- ▶ CPU scheduling is about selecting a task from the *Ready* queue to execute it on CPU, but when does the OS make such an action?

<sup>2</sup>Silberschatz, Galvin, and Gagne, *Operating system concepts*.

# Preemptive vs. Non-preemptive Scheduling

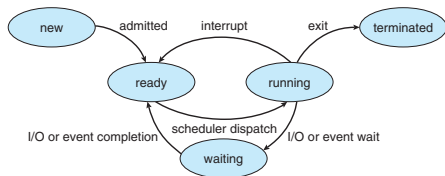
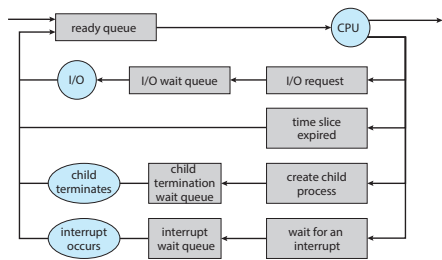


Figure: Process queues and transitions<sup>3</sup>.

Consider when a task goes into the *Ready* queue, or a task is off CPU,

1. Running  $\rightarrow$  Waiting
2. Running  $\rightarrow$  Ready
3. Waiting  $\rightarrow$  Ready
4. \*  $\rightarrow$  Terminated

<sup>3</sup>Silberschatz, Galvin, and Gagne, *Operating system concepts*.

# Preemptive vs. Non-preemptive Scheduling

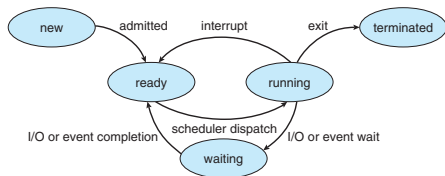
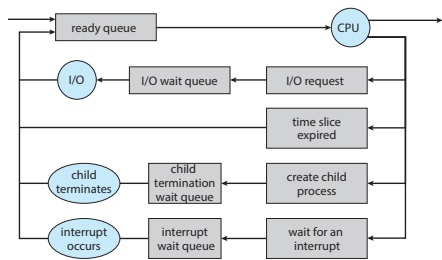


Figure: Process queues and transitions<sup>3</sup>.

Under 1 and 4, nonpreemptive or cooperative; otherwise, preemptive.

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2. Running  $\rightarrow$  Ready
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<sup>3</sup>Silberschatz, Galvin, and Gagne, *Operating system concepts*.

# Scheduling and Context Switch

1. CPU scheduler makes the decision and select a task from the *Ready* queue.
2. Dispatcher gives the control of the CPU to the selected task.
  - 2.1 Switching context from the active task to the selected CPU.
  - 2.2 Switching to user mode
  - 2.3 Jumping to the proper location in the selected task to resume that task



# Monitoring System Context Switching in Linux

```

1 -$ vmstat 1 3 # show vm statistics every 1 second for 3 times
2 procs -----memory----- ---swap-- -----io----- -system-- -----cpu-----
3  r  b   swpd   free   buff  cache   si   so    bi    bo    in   cs  us  sy  id  wa  st
4  0  0    1340 123188 112552 716608    0    0     0     2     4   10  0  0 100  0  0
5  0  0    1340 123188 112552 716640    0    0     0     0  576   74  0  0 100  0  0
6  0  0    1340 123188 112552 716640    0    0     0     0  576   68  0  0 100  0  0
7 -$

```

▶ `man vmstat`

`system`

`cs`: The number of context switches per second.

# Monitoring Process Context Switch in Linux

```

1  ~$ for p in /proc/[0-9]*; do \
2  > echo "Process ${p#/proc/}:"; \
3  > while read ln; do \
4  > echo -e "\t${ln}"; done <<< $(grep -E -o "\^.*_ctxt_switches.*$" ${p}/status); \
5  > done;
6  Process 1:
7     voluntary_ctxt_switches:          41215
8     nonvoluntary_ctxt_switches:      15741
9  Process 10:
10    voluntary_ctxt_switches:         26239510
11    nonvoluntary_ctxt_switches:      10
12    ...
13  Process 99:
14    voluntary_ctxt_switches:          4
15    nonvoluntary_ctxt_switches:      0
16  ~$

```

- ▶ A voluntary context switch occurs when a task has given up control of the CPU because it requires a resource that is currently unavailable (such as blocking for I/O.)
- ▶ A nonvoluntary context switch occurs when the CPU has been taken away from a task, such as when its time slice has expired or it has been preempted by an another task.

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## Not all processes are created equal

The design or selection of CPU scheduling algorithm depends on an observed property of processes:

- ▶ CPU burst and I/O burst cycle
- ▶ Distribution of CPU and I/O bursts
- ▶ CPU-bound processes
- ▶ I/O-bound processes

# Scheduling Criteria

Criteria from design or selection of CPU scheduling algorithm

- ▶ CPU utilization. % of time CPU being busy
- ▶ Throughput. # of tasks completed per time unit.
- ▶ Turnaround time. Interval from task submission to task completion.
- ▶ Waiting time. Total time a task spends (i.e., waits) in the ready queue.
- ▶ Response time. Interval from the submission of a request until the first response is produced.

# Optimizing for Scheduling Criteria

Maximize CPU utilization and throughput and to minimize turnaround time, waiting time, and response time.

- ▶ Consider min, max, average, variance ...
- ▶ Criteria may conflict with each other
- ▶ Batch systems vs. interactive systems vs. real-time systems vs. ...

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# Scheduling Algorithms

- ▶ First-Come, First-Served Scheduling
- ▶ Shortest-Job-First Scheduling
  - ▶ Shortest-Remaining-Time-First Scheduling
- ▶ Round-Robin Scheduling
- ▶ Priority Scheduling
  - ▶ Multilevel Queue Scheduling
  - ▶ Multilevel Feedback Queue Scheduling



# First-Come, First-Served Scheduling (FCFS)

Consider the *ready* queue with the following tasks,

Task	Burst Time
$P_1$	6
$P_2$	8
$P_3$	7
$P_4$	3

## Shortest-Job-First Scheduling (SJF)

Consider the *ready* queue with the following tasks,

Task	Burst Time
$P_1$	6
$P_2$	8
$P_3$	7
$P_4$	3

# Shortest-Remaining-Time-First Scheduling (SRTF)

Consider the following tasks that arrive in the *ready* queue,

Task	Arrival Time	Burst Time
$P_1$	0	6
$P_2$	2	8
$P_3$	3	7
$P_4$	6	3

## Estimating CPU Burst Time

- ▶ Do we know the CPU burst times at the time when we invoke the CPU scheduling algorithm?

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- ▶ How do we predict CPU burst times?  
The next CPU burst is generally predicted as an *exponential average* of the measured lengths of previous CPU bursts, commonly,

$$\tau_{n+1} = \alpha t_n + (1 - \alpha)\tau_n \quad (1)$$

where

$t_n$ : the n-th CPU burst that the OS records

$\tau_{n+1}$ : the next (i.e., n+1) predicted value of the CPU burst

$\alpha$ :  $\alpha \in [0, 1]$ , an aging exponent that determines the effect of history of CPU bursts.

## Round-Robin Scheduling (RR)

Assume time quantum = 2 and consider the *ready* queue with the following tasks,

Task	Burst Time
$P_1$	6
$P_2$	8
$P_3$	7
$P_4$	3

# Priority Scheduling

Consider the *ready* queue with the following tasks where 1 means the highest priority, and 3 lowest,

Task	Burst Time	Priority
$P_1$	6	3
$P_2$	8	1
$P_3$	7	2
$P_4$	3	1



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# Thread Scheduling

- ▶ Kernel and user threads
- ▶ Contention scope

# Multiprocessor Scheduling

- ▶ Multiprocess architecture, multicore CPUs vs. multithreaded cores vs. NUMA systems vs. Heterogeneous multiprocessing
- ▶ Common ready queue vs. per-core ready queue
- ▶ Load balancing
- ▶ Processor affinity and cache

# References I



Silberschatz, Abraham, Peter B. Galvin, and Greg Gagne. *Operating system concepts*. 10th edition. John Wiley & Sons, 2018.