Simple Internetworking: Global Addressing and Datagram Forwarding

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Outline

- Topic: internetworking
	- Case study: Internet Protocol (IP) Suite
- Simple interworking
	- Overview of internet and the **I**nternet
	- Global addressing scheme
	- **Best effort service model and datagram forwarding**
	- Packet fragmentation and assembly
	- Address translation
	- Host configuration
	- Error reporting

IP Service Model

- Packet Delivery Model
	- Connectionless model for data delivery
	- Best-effort delivery (unreliable service)
		- packets may be lost
		- packets may be delivered out of order
		- duplicate copies of a packet may be delivered
		- packets may be delayed for a long time
	- Why motivated this design requirement (policy)?
- Mechanism (solution) to support this design requirement (policy)
	- Datagram forwarding between networks (application of datagram switching)
	- Two problems:
		- Identify hosts
		- Identify networks

IP Address Space

- Global unicast addresses
- Multicast addresses
- Broadcast addresses
- Private addresses
- Link-local addresses
- And more (such as for testing and documentation)

IPv4 Address Space

- 32-bit long
- 2³² Addresses
	- Global unicast addresses
	- Multicast addresses
	- Broadcast addresses
	- Private addresses (private unicast addresses)
	- Link-local addresses

• …

Global Unicast IPv4 Addresses

- Binary representation begins with 0, 10, 110
- Globally unique
- Human-readable form: IPv4 numbers-and-dots notation
	- 146.245.201.50
- hierarchical: network + host
	- An unicast IPv4 address specifies both network and the host on the network

IPv4 Multicast Address

• Addresses starting with 1110

Multicast Address 1 1 1 0 Group Address

IPv4 Broadcast Address

- IP unicast address: network number + host number
- Setting all the host bits to 1

Network 11111...11111

• We shall discuss methods to divide the bits into network number and host number

IPv4 Address Spaces for Private Networks

- See [RFC 1918](https://tools.ietf.org/html/rfc1918)
- Private networks
	- 24-bit block 10.0.0.0–10.255.255.255
		- First 8 bits: 0000 1010
	- 20-bit block 172.16.0.0–172.31.255.255
		- First 12 bits: 1010 1100 0001
	- 16-bit block 192.168.0.0–192.168.255.255
		- First 16bits: 1100 0000 1010 1000

Link Local IPv4 Address

■ See [RFC 3927](https://tools.ietf.org/html/rfc3927)

□ Link-Local IPv4 Address

- 16-bit block 169.254.0.0–169.254.255.255
- First 16 bits: 1010 1001 1111 1110

Exercise 1

- Find out IPv4 addresses of following hosts (with the given domain names) and indicate the class to which the IP addresses belong
	- www.brooklyn.cuny.edu
	- www.sme.sk
	- www.google.com
- Convert the first number (from left) to a binary number, then take a look at the 1^{st} , and/or 2^{nd} , and/or 3^{rd} bit
- Are the addresses global unicast addresses?
- Remark
	- There are many ways to find out the IP address of a host given a domain name
		- Example: nslookup (which works on most platforms including Windows, Unix/Linux, and Mac OS X) and dig (works on Linux)

Summary and Questions?

- IP service model
- IPv4 addresses
	- Unicast addresses
	- Multicast addresses
	- Broadcast addresses
	- Private addresses
	- Link-local addresses
- Look-up global unicast address assignment for some hosts with domain names

IP Datagram Forwarding

- Set-up:
	- Every datagram contains source and destination's addresses
		- Each address: network number + host number
	- Each network has network gateways/routers
	- Each host and router maintains a forwarding table
		- each host/router can have a default router
	- forwarding table maps network number into next hop
- Forwarding algorithm
	- if source and destination are on the same network, then forward to the destination host
	- If source and destination are not connected to the same network, then forward to some router (representing a network)

Forwarding Algorithm

if (NetworkNum of destination = NetworkNum of one of my interfaces) **then**

deliver packet to destination over that interface

else

if (NetworkNum of destination is in my forwarding table) **then**

deliver packet to NextHop router

else

deliver packet to default router

Forwarding Algorithm Simplified

- For a host with only one interface and only a default router in its forwarding table, this simplifies to
- **if** (NetworkNum of destination = my NetworkNum)**then** deliver packet to destination directly

else

deliver packet to default router

Forwarding Table: Example

• Forwarding table at router R2 that has two interfaces 0 and 1

Exercise 2

• Construct forwarding tables for routers R1 and R3. Interfaces of routers are marked

Summary and Questions?

- IP Datagram Forwarding
- Forwarding table
- Forwarding algorithm
- But how do we divide an address into a network number and a host number?

Network + host numbers?

- We will need to divide a unicast address into two parts: network number + host numbers
	- (legacy) IPv4 address classes
	- (legacy) subnetting/supernetting
	- CIDR addressing
- Forwarding algorithms will need to be adjusted accordingly **NovtHon NaturarkNum**

IPv4 Unicast Address Classes

- To express network and host
	- Divide into classes (legacy)

Subnetting

- Subnetting:
	- divide large network into smaller networks using a network mask
	- merge smaller networks into a large network using a network mask

Subnetting: Example

Subnetting: Example: Forwarding Table

• Forwarding Table at Router R1

Subnetting: Discussion

- Would use a default router if nothing matches
- Subnet masks do not have to align with a byte boundary
- Subnet masks need **not** to be contiguous 1's
	- 255.255.1.0 is OK
		- 11111111 11111111 00000001 00000000
		- What is subnet number of IP address 128.96.34.1? 10000000 01100000 00100010 00000000 & 11111111 11111111 00000001 00000000 \rightarrow 10000000 01100000 00000000 00000000 \rightarrow 128.96.0.0 \rightarrow can not directly tell from the IP address
	- In practice, use contiguous 1's
- Multiple subnets can be on a single physical network
- Subnets not visible from the rest of the Internet

Subnetting: Discussion

• How do you tell whether an IP address is on a given subnet?

Subnet mask: 255.255.255.0 Subnet number: 128.96.33.0

Forwarding Algorithm

```
D = destination IP address
for each entry < SubnetNum, SubnetMask, NextHop> 
 in the forwarding table
   D1 = SubnetMask & D
   if D1 = SubnetNum
      if NextHop is an interface
          deliver datagram directly to destination
      else
          deliver datagram to NextHop (a router)
```
Exercise 3

- The router in an IPv4 network uses subnetting. State to what next hop the IP packets addressed to each of the following destinations will be delivered
	- (a) $128.96.171.92$
	- (b) 128.96.167.151
	- (c) 128.96.163.151
	- (d) 128.96.169.192
	- (e) 128.96.165.121

Classless Addressing

• Represent network number with a single pair

<length, value>

• All routers must understand CIDR addressing

Classless Addressing: Notation

- Also called CIDR addressing (**C**lassless **I**ntra**d**omain **R**outing addressing)
- Convention
	- Place a /X after the prefix where X is the prefix length in bits
- Example
	- 20-bit prefix for all the networks 192.4.16 through 192.4.31: 192.4.16/20
	- A single class C network number, which is 24 bits long: 192.4.16/24

Handling Classless Addresses

- How do the routing protocols handle this classless addresses
	- It must understand that the network number may be of any length
	- Requires to hand out blocks of addresses that share a common prefix
	- Revising IP forwarding algorithm

IP Forwarding Revisited

- Original *assumptions* in IP forwarding mechanism
	- It can find the network number from destination IP address in a packet
	- Then look up that number in the forwarding table
- Need to *change* this assumption in case of CIDR
	- How? (next slide)

IP Forwarding in CIDR

- Prefixes may be of any length
	- Prefixes in the forwarding tables may overlap
- Some addresses may match more than one prefix
	- Example
		- Both 171.69/16 and 171.69.10/24 may coexist in the forwarding table of a single router
		- A packet destined to 171.69.10.5 clearly matches both prefixes.
	- Principle of "longest match"
		- A packet destined to 171.69.10.5 matches prefix 171.69.10/24
		- A packet destined to 171.69.20.5 matches 171.69/16

Exercise 4

• State to what next hop the IP packets addressed to each of the following destinations will be delivered

(a) $C4.4B.31.2E$ (b) C4.5E.05.09

- (c) $C4.4D.31.2E$
- (d) C_{4.5}E_{.03.87}
- (e) C4.5E.7F.12
- (f) C4.5E.D1.02

Benefits of Classless Addressing

- Addressing two scaling challenges
	- Potential exhaustion of the 32-bit address space
	- The growth of backbone routing table as more and more network numbers need to be stored in them (route aggregation)
- Address assignment efficiency
	- Arises because of the IP address structure with class A, B, and C addresses
	- Forces us to hand out network address space in fixed-size chunks of three very different sizes
		- A network with two hosts needs a class C address:
			- Address assignment efficiency = 2/255 = 0.78
		- A network with 256 hosts needs a class B address
			- Address assignment efficiency = 256/65535 = 0.39

Classless Addressing: Address Assignment: Example

- Consider an organization has 16 class C network numbers.
- Instead of handing out 16 addresses at random, hand out a block of contiguous class C addresses
- Suppose we assign the class C network numbers from 192.4.16 through 192.4.31
- Observe that top 20 bits of all the addresses in this range are the same (11000000 00000100 0001)
	- We have created a 20-bit network number (which is in between class B network number and class C number)
- Requires to hand out blocks of class C addresses that share a common prefix (sometimes, called supnetting)

Route Aggregation with CIDR: Example

• Route aggregation with CIDR

Summary

- Global addressing scheme
- Best effort service model and datagram forwarding
	- Defining a network
		- Legacy IPv4 address classes
		- Subnetting (supernetting)
		- Classless addressing (CIDR addressing)