Direct Link Networks: Encoding and Framing

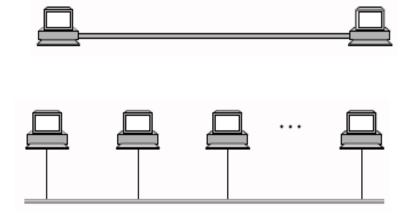
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Direct Link Networks

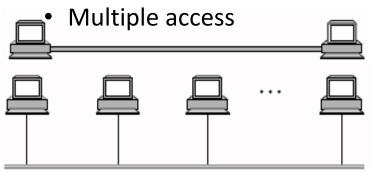
- Types of Networks
 - Point-to-point
 - Multiple access

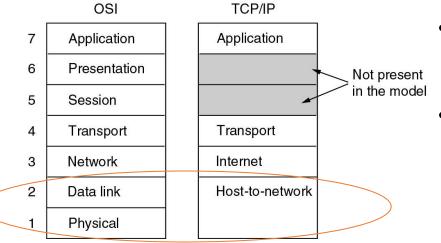


 What problems do we need to solve to build a direct link network?

Direct Link Networks

- Types of Networks
 - Point-to-point

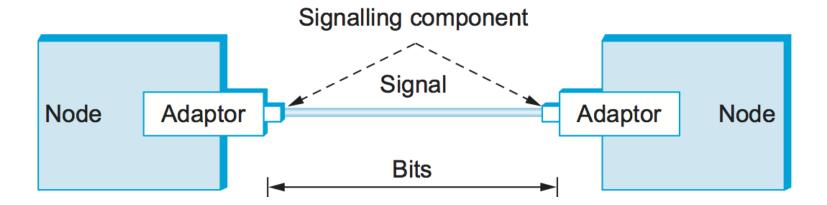




- Encoding
 - Encoding bits onto transmission medium
- Framing
 - Delineating sequence of bits into messages
- Error detection
 - Detecting errors and acting on them
- Reliable delivery
 - Making links appear reliable despite errors
- Media access control
 - Mediating access to shared link

Encoding

- Encoding
 - Encode bits (binary data) into the signals
 - Modulation is not our focus
 - Assume working with two discrete signals: high and low

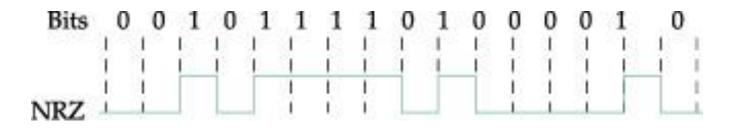


Comparing Encoding Schemes

- Nonreturn-to-zero (NRZ)
- Nonreturn-to-zero-inverted (NRZI)
- Manchester
- 4B/5B

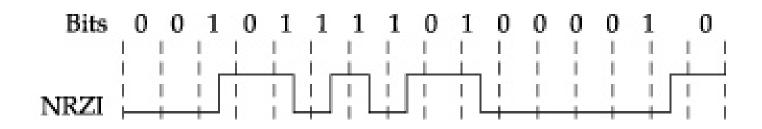
Non-Return-to-Zero (NRZ)

- Method
 - Low \rightarrow 0
 - High $\rightarrow 1$
- Shortcoming
 - Long strings of 1s or 0s
 - Baseline wander
 - Difficult to recover clock



Non-Return-to-Zero-Inverted (NRZI)

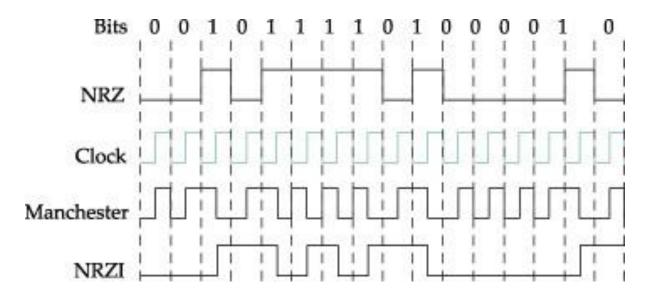
- Method
 - Signal transition \rightarrow 1
 - No transition \rightarrow 0
- Solve the problem caused by consecutive 1's
 - But, the problem caused by consecutive 0's remains



Manchester

- Method
 - 0 \rightarrow low-to-high transition

- $1 \rightarrow$ high-to-low transition
- That is, NRZ signal \oplus Clock signal
- Solve the problems caused by both consecutive 1's and 0's
- New problem
 - 50% efficient: Two samples per clock cycle to detect a transition, bit rate is 50% of baud rate (rate of signal changes)



Exercise

- Encode bit sequence 01101
 - using the NRZ, NRZI, Manchester encodings
 - indicating bits and by drawing clock, NRZ, NRZI, and Manchester signals
- (example in the previous slide)

4B/5B

- Motivations
 - To cope with baseline wander
 - To ease clock recovery
 - To increase encoding efficiency (bit rate/baud rate)

4B/5B

• Break long strings of repeated 0's and 1's

- before transmission: 4 bits \rightarrow 5 bits
 - Every 4 bits of actual data are encoded in a 5-bit code
 - 4 bits data symbols: 2⁴ = 16
 - 5 bits codes: 2⁵ = 32
 - No more than 1 leading 0's in codes and no more than 2 trailing 0's
 - No pair of 5-bit codes results in more than 3 consecutive 0's
 - 11111: line is idle 01101: control symbols
- Transmit codes using NRZI: 80% efficiency

4B/5B Encoding

| 4-Bit Data Symbol | 5-Bit Code | 4-Bit Data Symbol | 5-Bit Code |
|-------------------|------------|-------------------|------------|
| 0000 | 11110 | 1000 | 10010 |
| 0001 | 01001 | 1001 | 10011 |
| 0010 | 10100 | 1010 | 10110 |
| 0011 | 10101 | 1011 | 10111 |
| 0100 | 01010 | 1100 | 11010 |
| 0101 | 01011 | 1101 | 11011 |
| 0110 | 01110 | 1110 | 11100 |
| 0111 | 01111 | 1111 | 11101 |

Exercise

 Encode bit sequence 01011000 using 4B/5B encoding by showing the conversion between 4-bit sequences of data and their 5-bit sequences of codes and drawing clock and 4B/5B signals

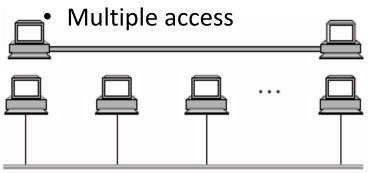
| 4-Bit Data Symbol | 5-Bit Code | 4-B | it Data Symbol | 5-Bit Code |
|-------------------|------------|------|----------------|------------|
| 0000 | 11110 | 100 | 0 | 10010 |
| 0001 | 01001 | 100 | 1 | 10011 |
| 0010 | 10100 | 101 | 0 | 10110 |
| 0011 | 10101 | 101 | 1 | 10111 |
| 0100 | 01010 | 110 | 0 | 11010 |
| 0101 | 01011 | 110 | 1 | 11011 |
| 0110 | 01110 | 1110 | D | 11100 |
| 0111 | 01111 | 111 | 1 | 11101 |

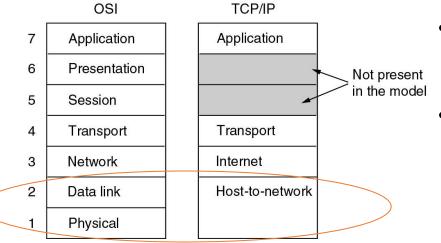
Summary: Encoding

- Encoding problem
- Solutions
 - NRZ
 - NRZI
 - Manchester
 - 4B/5B
- Advantages/disadvantages?

Direct Link Networks

- Types of Networks
 - Point-to-point





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Framing

- Packet-switched networks → block of data are exchanged between nodes
 - Breaking bits into frames (packets)
- Key issue: identity where a frame begins and the ends
- Why do we need it?
- To discuss in the context of point-to-point links
 - Byte-oriented framing: frame as a collection of bytes
 - Bit-oriented framing: frame as a collection of bits
 - Clock-based framing (vs. sentinel-based approaches)
- Framing is a fundamental problem that must be addressed in multiaccess networks as well.

Byte-Oriented Framing

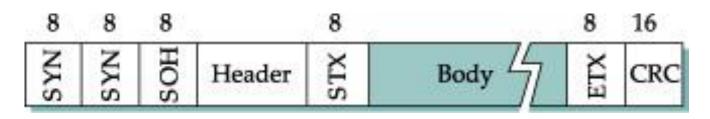
- Each frame as a collection of bytes.
- Examples
 - Binary Synchronous Communication (BISYNC) protocol
 - Digital Data Communication Message Protocol (DDCMP)
 - Point-to-Point Protocol (PPP)
 - widely used today

Sentinel vs Length

- Using sentinels
 - Use special characters to indicate where frames start and end.
 - e.g., BISYNC
 - SYN (synchronization), STX (start of text), ETX (end of text)
 - Need "escaping"/character stuffing : special characters might appear in the data portion of the frame (DLE, or data link escape)
- Using a length value
 - include the number of bytes in the frame at the beginning of the frame, i.e., in the frame header.
 - Example: DDCMP
 - A transmission error could corrupt the count field, in which case the end of the frame would not be correctly detected.
 - framing error

BISYNC

- Uses special characters known as *sentinel* characters to indicate whether frames begins and ends
- The beginning of a frame: SYN
- The data portion of the frame is contained between STX and ETX
- The start of header: SOH



BISYNC Character Stuffing

- Problem: ETX may appear in the data
- Solution: escaping
 - ETX \rightarrow DLE ETX
 - \cdot DLE \rightarrow DLE DLE

Point-to-Point Protocol (PPP)

- Commonly used to carry Internet Protocol packets over various sorts of point-to-point links
- Uses sentinels and character stuffing.

| 8 | 8 | 8 | 16 | | | 16 | 8 |
|------|---------|---------|----------|---------|---|----------|------|
| Flag | Address | Control | Protocol | Payload | 7 | Checksum | Flag |

Illustrating Frame Format

- A frame is illustrated as a sequence of labeled fields
 - Above each field is a number indicating the length of that field in bits.
 - Note that the packets are transmitted beginning with the leftmost field.
 - Often we can find more detailed explanation of each field

| 8 | 8 | 8 | 16 | | 16 | 8 |
|------|---------|---------|----------|---------|----------|------|
| Flag | Address | Control | Protocol | Payload | Checksum | Flag |

Illustrating Frame Format: Fields

| 8 | 8 | 8 | 16 | | | 16 | 8 |
|------|---------|---------|----------|---------|---|----------|------|
| Flag | Address | Control | Protocol | Payload | 7 | Checksum | Flag |

- The special start-of-text character, denoted as the Flag field is 01111110.
- The Address and Control fields usually contain default values.
- The (Protocol) field is used for demultiplexing; it identifies the high-level protocol, such as IP.
- The frame payload size can be negotiated, but it is 1500 bytes by default.
- The Checksum field is either 2 (by default) or 4 bytes long.
- The PPP frame format is unusual in that several of the field sizes are negotiated rather than fixed.
 - This negotiation is conducted by a protocol called the Link Control Protocol (LCP). PPP and LCP work in tandem
 - LCP sends control messages encapsulated in PPP frames

Exercise

- Use the sentinels in the ASCII table.
 - (Linux manual page) man ascii
- DATA = 1A E2 02 2A 16 10 20. What would be the bytes in the *body* of the frame using BISYNC?
- In the *body* of a frame using BISYNC, is it possible to see the following byte sequence and why?

4A 10 51 6B

Bit-Oriented Framing

- View a frame as a collection of bits
 - The Synchronous Data Link Control (SDLC) protocol
 - High-Level Data Link Control (HDLC) protocol

HDLC

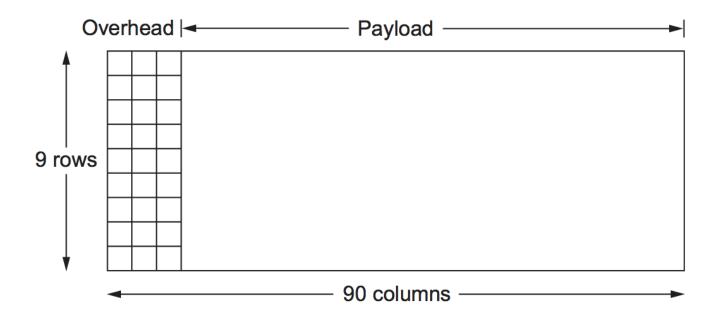


- Beginning and ending sequence: 01111110
- 01111110 may appear in the body of the frame → bit stuffing
 - Sender: any time five consecutive 1s have been transmitted from the body of the, the sender inserts a 0 before transmitting the next bit.
 - 11111 \rightarrow 111110 \rightarrow no 6 1's in a row in the body
 - Receiver: three scenarios
 - Bit stuffed; end of frame marker; error

Clock-based Framing

- Versus sentinel-based approaches
 - Due to bit stuffing or character stuffing, it is in fact not possible to make all frames exactly the same size,
- Clock-based framing
 - Frame can be of fixed length
- Example: framing in SONET
 - Popular in network backbone
 - Has a close tie with telephony systems
 - Specification takes entire book
 - Frame is of fixed length
 - 125 µs for STS-1 (51.84 Mbps)
 - Q: How big is a Mega (sidebar in page 45)
 - 125 µs = ? Bytes
 - Detect the frame header at each interval of frame size

SONET STS-1 Frame



- Each frame is 9 x 90 = 810 bytes long
- The first 2 bytes of the frame contain a special bit pattern
- When the special pattern turns up in the right place enough times, the receiver concludes that it is in sync and can then interpret the frame correctly.

Summary

- Bits/Bytes → frames
 - Framing
- Q: What if the link is not error free? In other words, what if a frame is corrupted?
 - Error detection
 - Reliable transmission